

# Giving State to the Stateless: Augmenting Trustworthy Computation with Ledgers

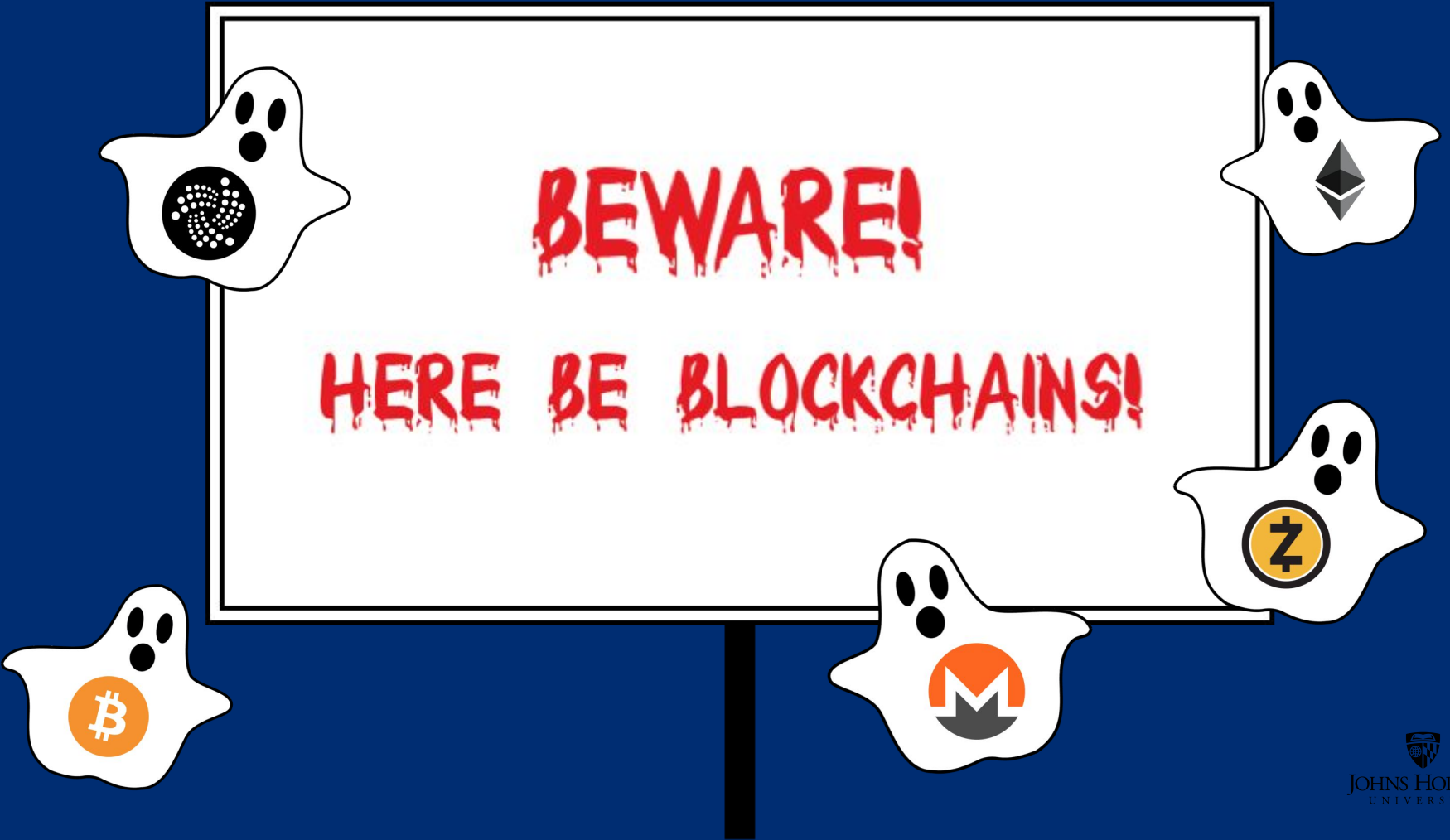
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*Gabriel Kaptchuk*, Matthew Green, and Ian Miers

**BEWARE!**

**HERE BE BLOCKCHAINS!**





**BEWARE!**

**HERE BE SGX!**



Now that we have SGX, all security problems are trivial. Security research is officially over!

HERE BE





**BEWARE!**

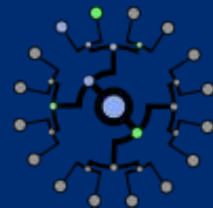
**HERE BE SGX!**



# Why care about the bogeymen?

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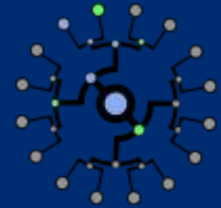
- Ledgers exist *in practice* and they aren't going away
  - Blockchains
  - Google Certificate Transparency Log
  
- Trusted Execution Environments exist *in practice*



# Why care about the bogeymen?

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- Ledgers exist *in practice* and they aren't going away
  - Blockchains
  - Google Certificate Transparency Log



- Trusted Execution Environments exist *in practice* ... kinda?

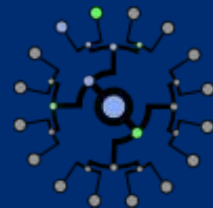


# Why care about the bogeymen?

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- Ledgers exist *in practice* and they aren't going away

- Blockchains
- Google Certificate Transparency Log



- Trusted Execution Environments exist *in practice* ... kinda?

- Intel SGX and ARM Trustzone
- Software only obfuscation
- FPGA style hardware with burned keys

**How can TEE's augment ledgers?**  
**VS**  
**How can ledgers augment TEE's?**

# Rewind Protection

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Is the password "1234"?



Nope! You have 9 more attempts!



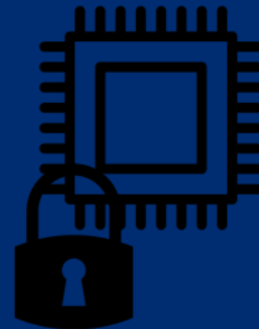
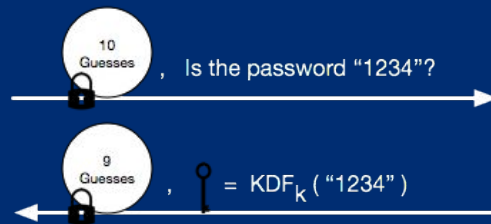
# Rewind Protection



Is the password "1234"?



Decryption Failure! 9 more attempts!



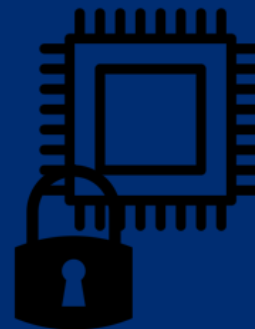
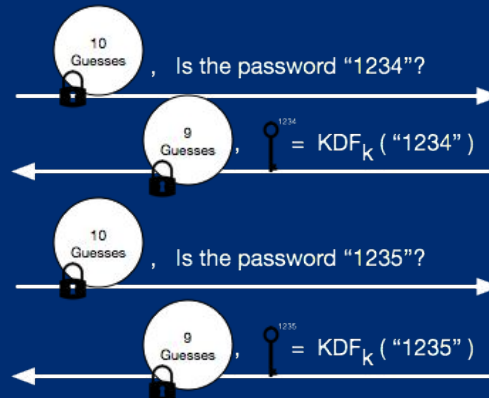
# Rewind Protection



Is the password "1234" or "1235"?



Decryption Failure! 9 more attempts!



# Rewind Protection

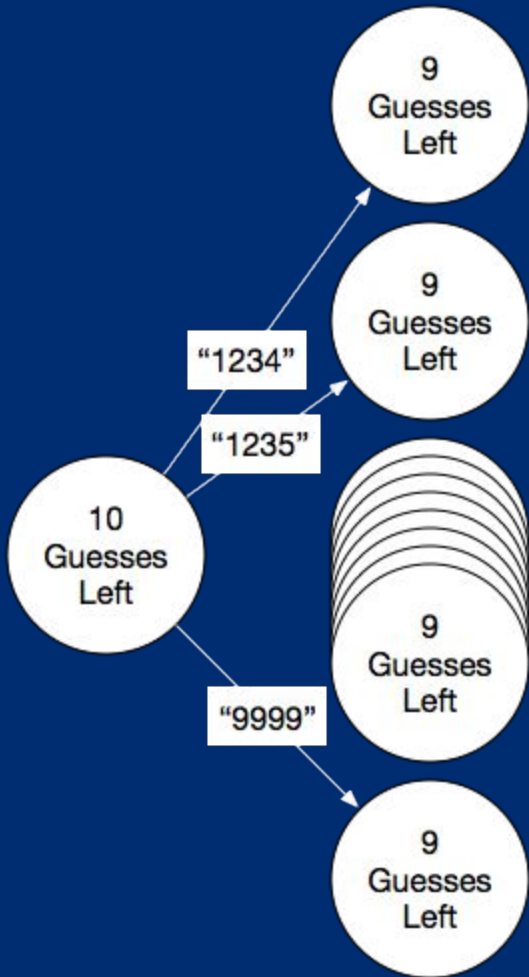
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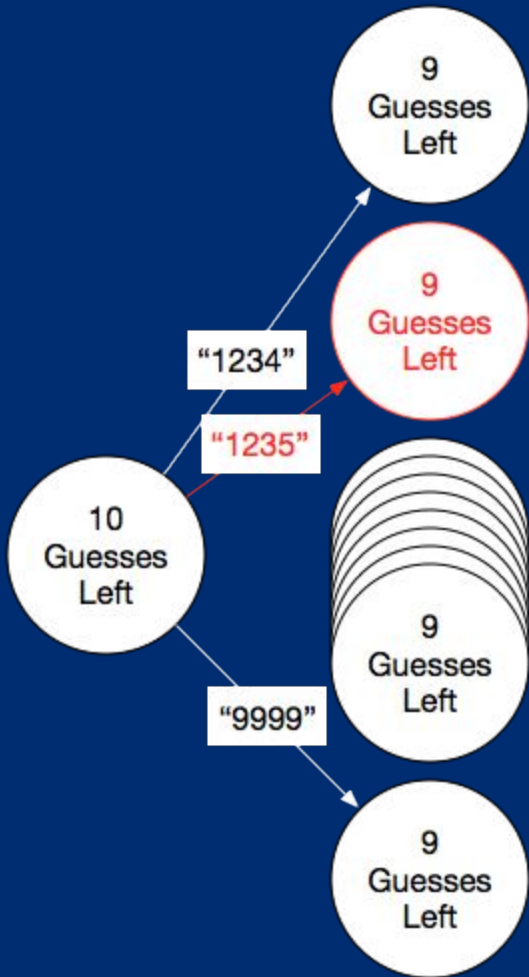
- Hardware based TEE's require NVRAM for protection
  - Scale poorly, expensive, and require special considerations for power fluctuations
  - Prior Work: Memoir [PLDMM11]
  
- Software only obfuscation can't get hardware-back protections
  - Prior Work: Goyal and Goyal [GoyGoy17] get one time programs from Ledgers + Obfuscation
  
- This problem is *real*

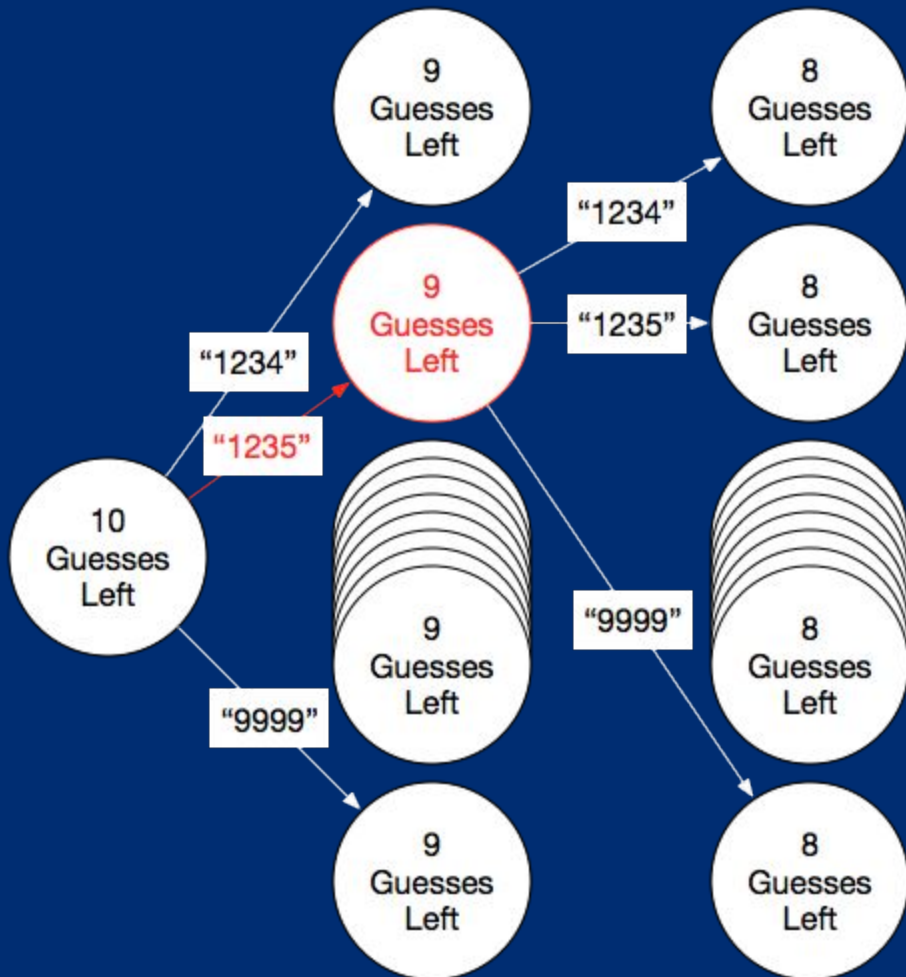


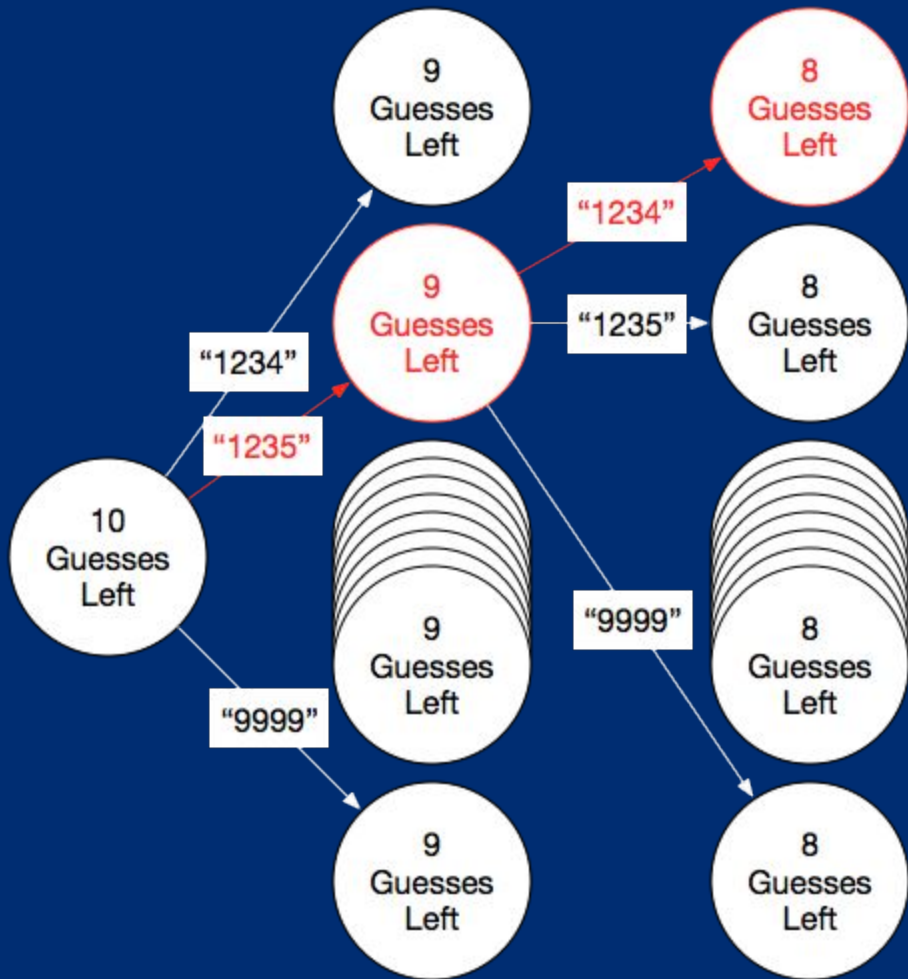
10  
Guesses  
Left





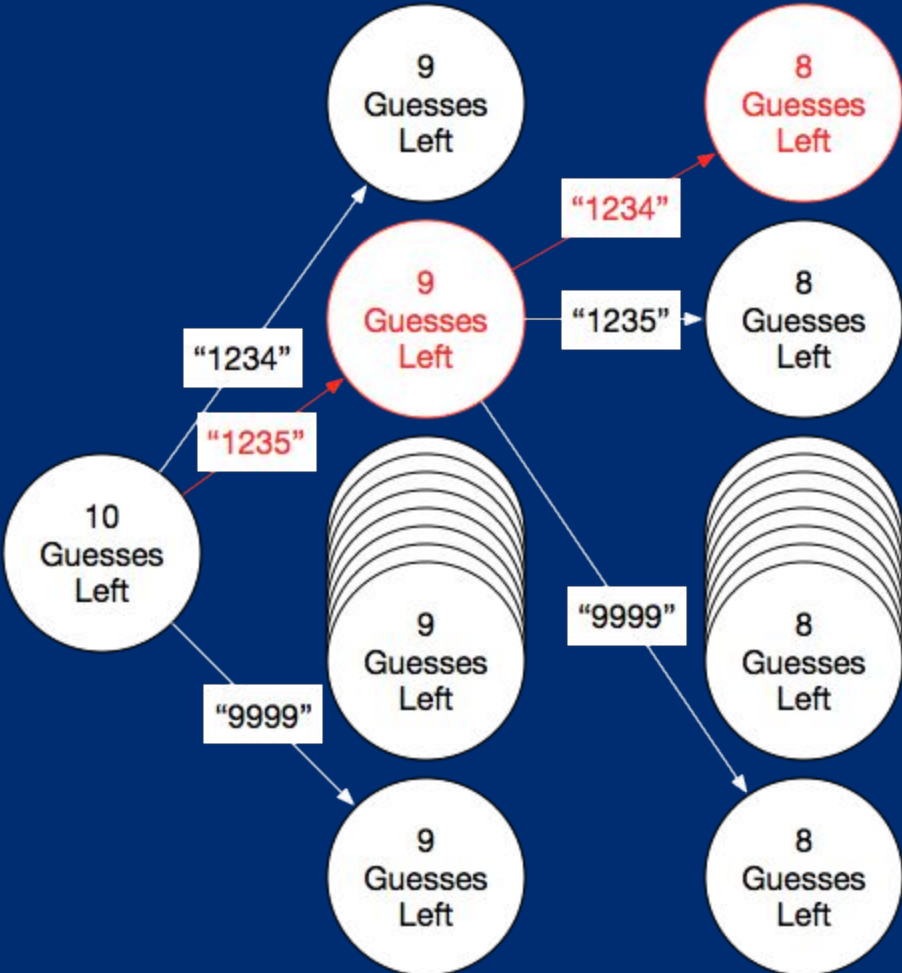


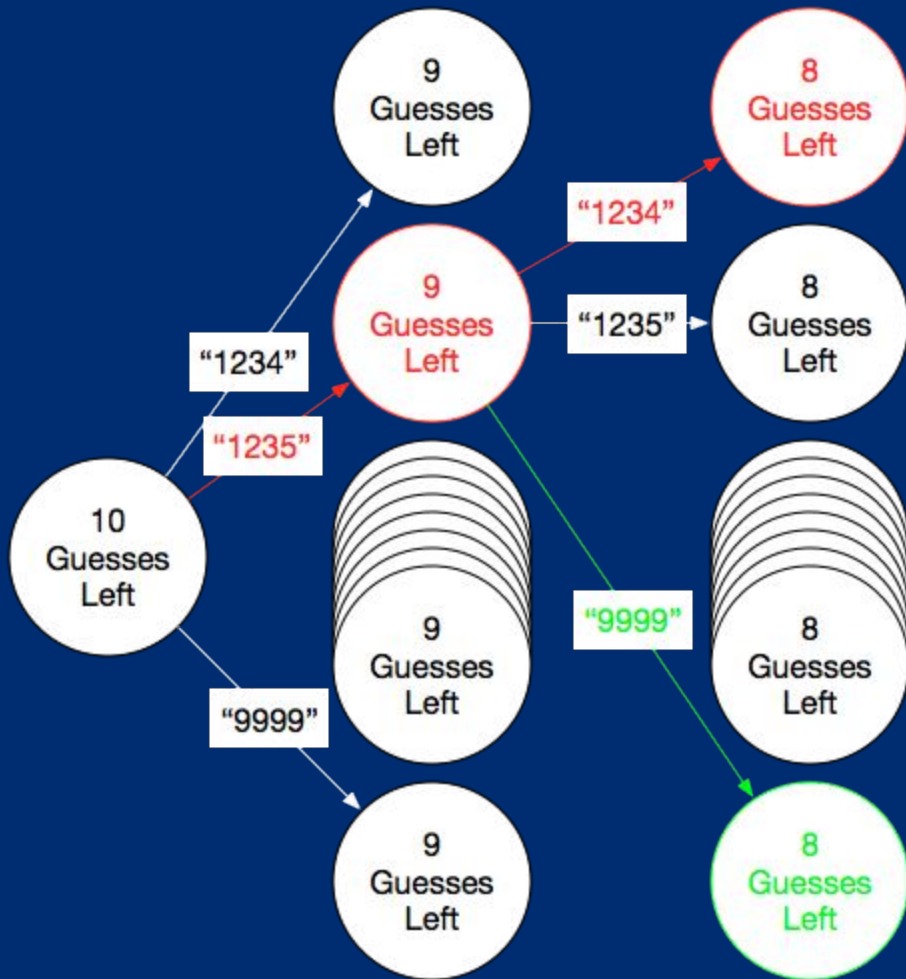




# Repeated Execution

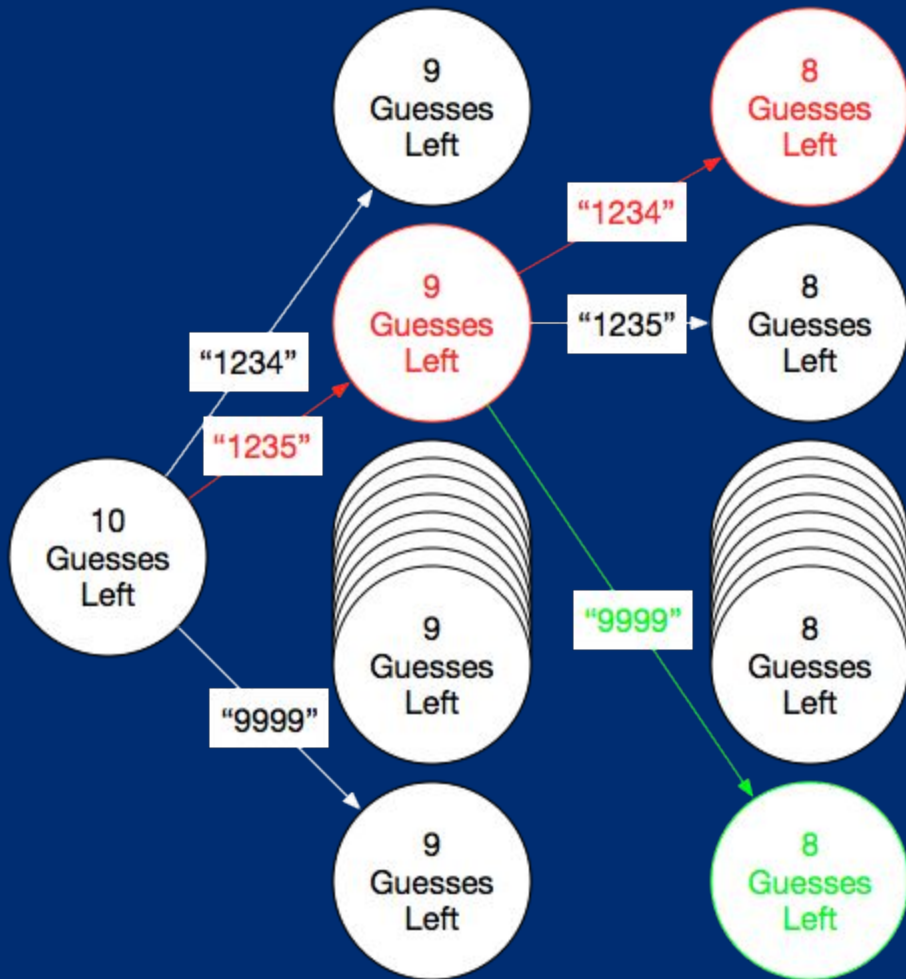
- Re-execution of a path doesn't cause a vulnerability
  - Derive the same key repeatedly
  - Starting again generates new master key





## Repeated Execution

- Re-execution of a path doesn't cause a vulnerability
  - Derive the same key repeatedly
  - Starting again generates new master key
- Forking is dangerous
  - Running new inputs on old state
  - Running old steps with new randomness



## Repeated Execution

- Re-execution of a path doesn't cause a vulnerability
  - Derive the same key repeatedly
  - Starting again generates new master key
- Forking is dangerous
  - Running new inputs on old state
  - Running old steps with new randomness
- Strategy: bind program execution to something linear

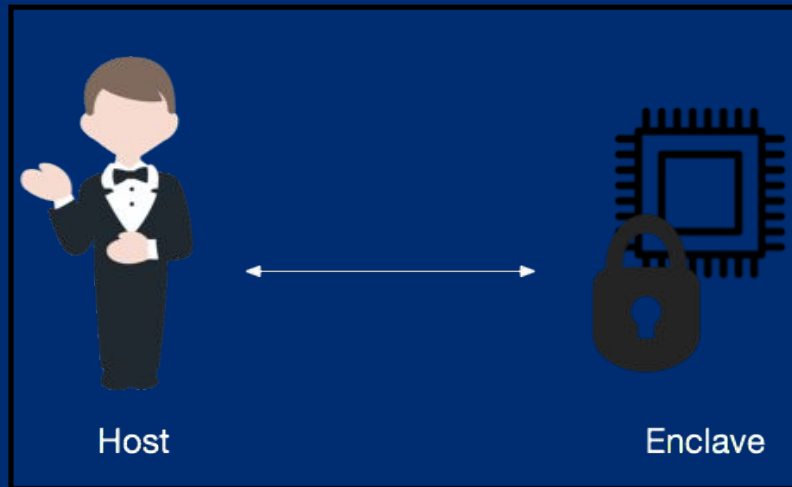
# Model

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Ledger



User



Host

Enclave



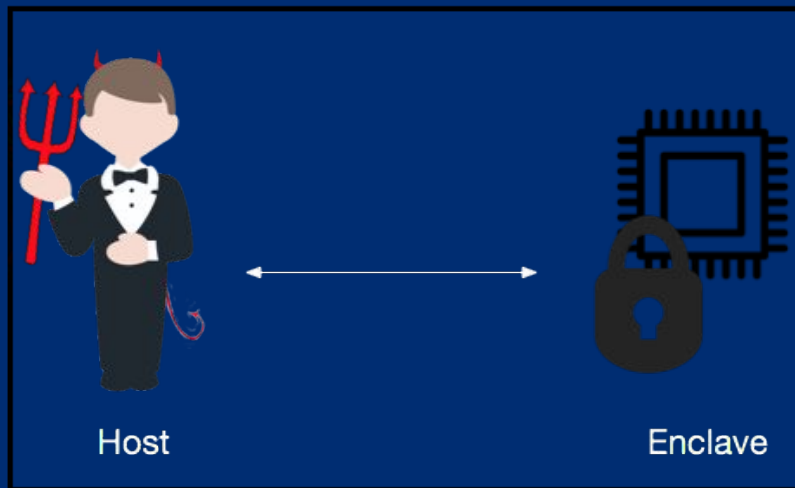
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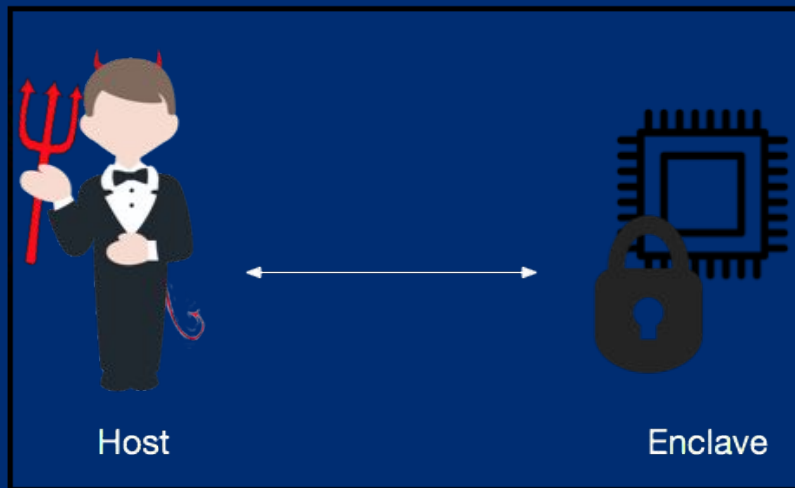
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# Model

Ledger



User



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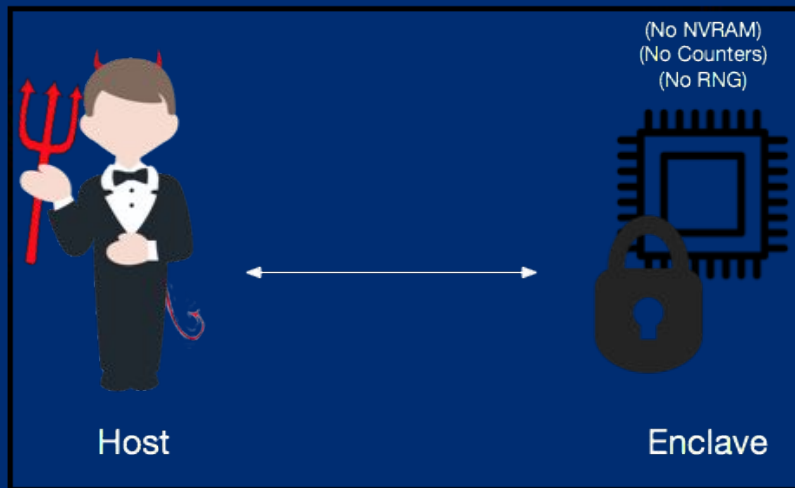
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Ledger



User



Host

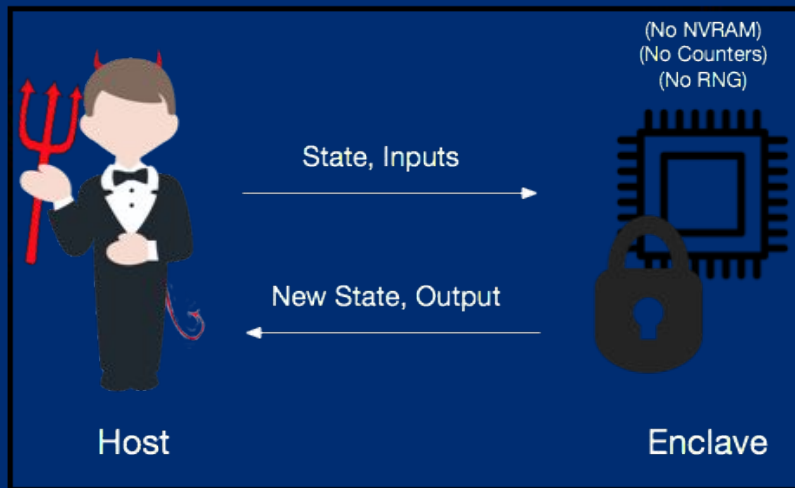
Enclave

# Model

Ledger



User



Host

Enclave

# Model

Ledger

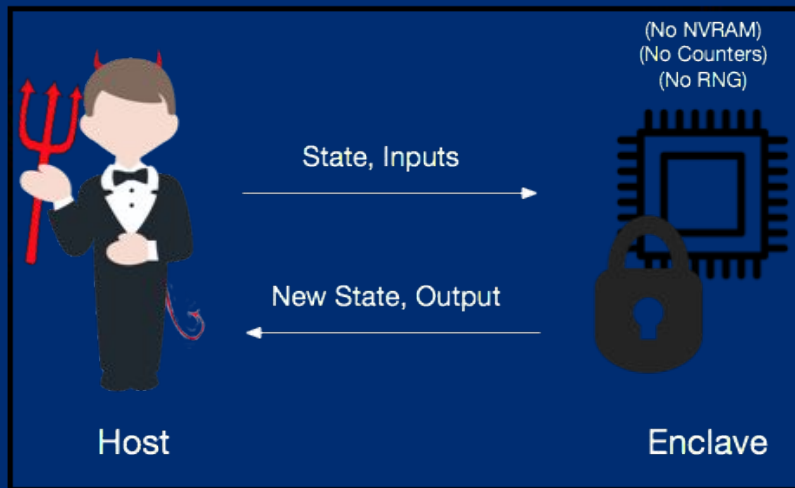


Data

$\sigma$



User



Host

Enclave

# Model

Ledger



Data

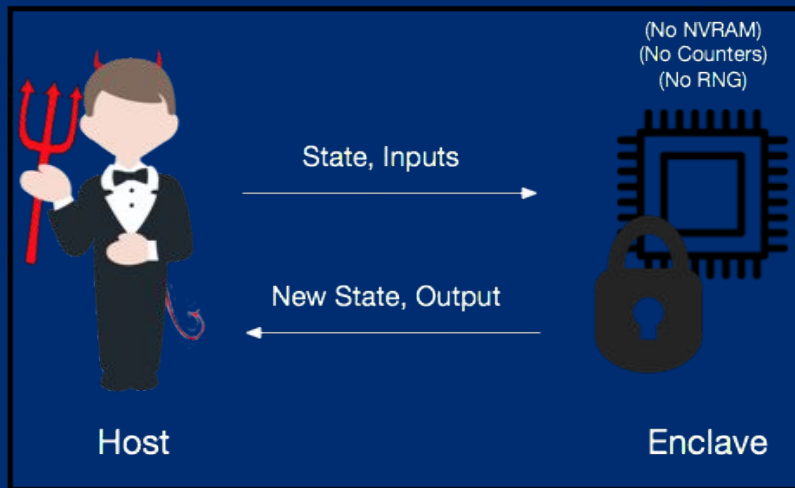
$\sigma$



User

Inputs

Output



State, Inputs

New State, Output

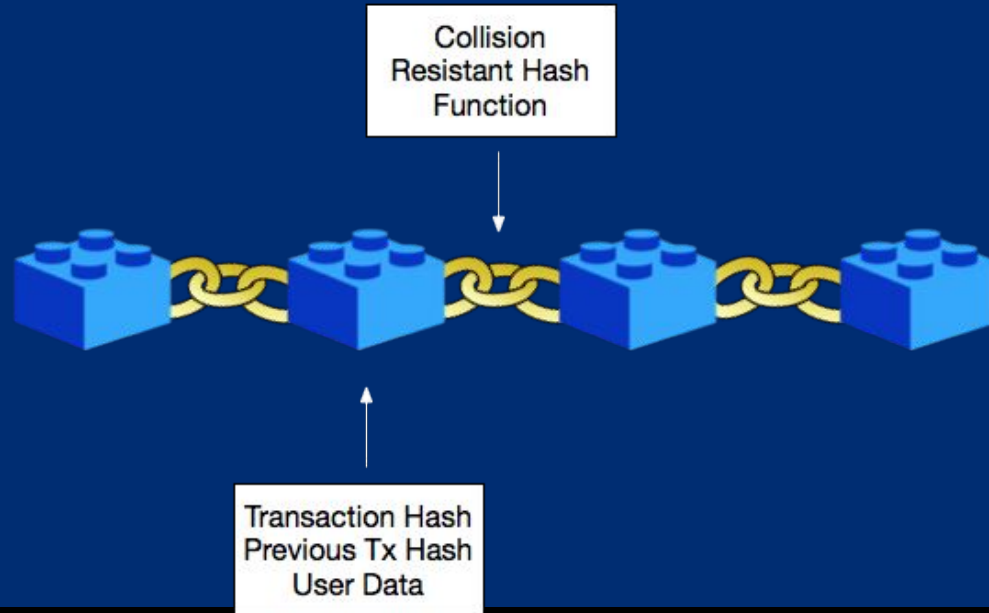
Host

Enclave

# Ledger Requirements

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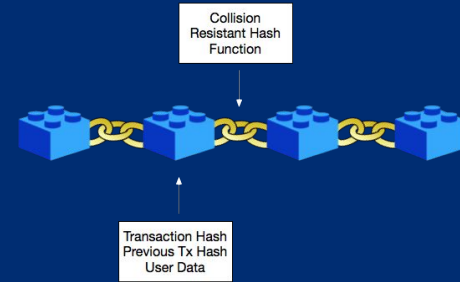
- Creates hash chains of transactions
  - Similar to transaction in bitcoin, ethereum, etc...



# Ledger Requirements

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- Publicly verifiable proof of publication and public access
  - Digital signatures for computational security
  - Proof of work for economic security

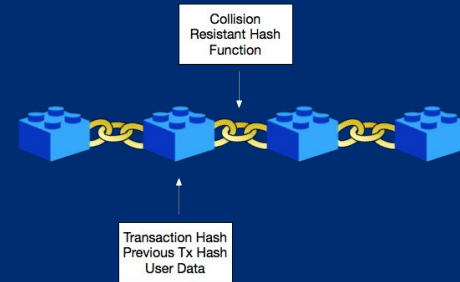




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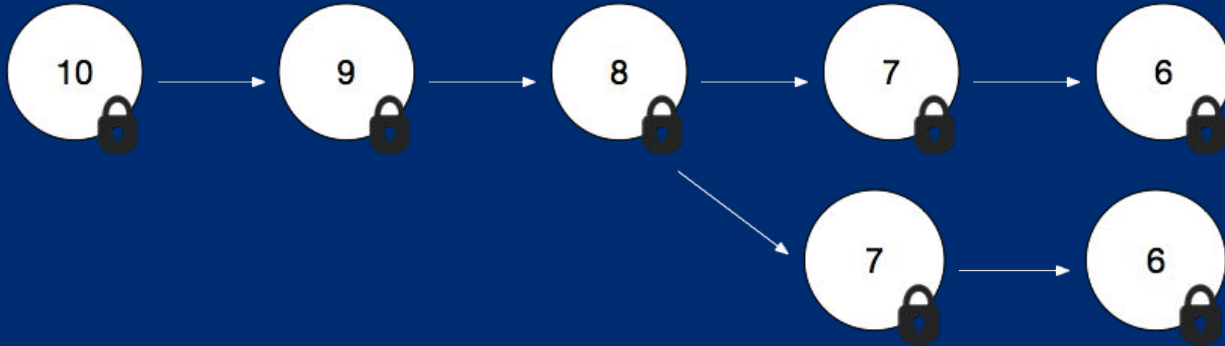
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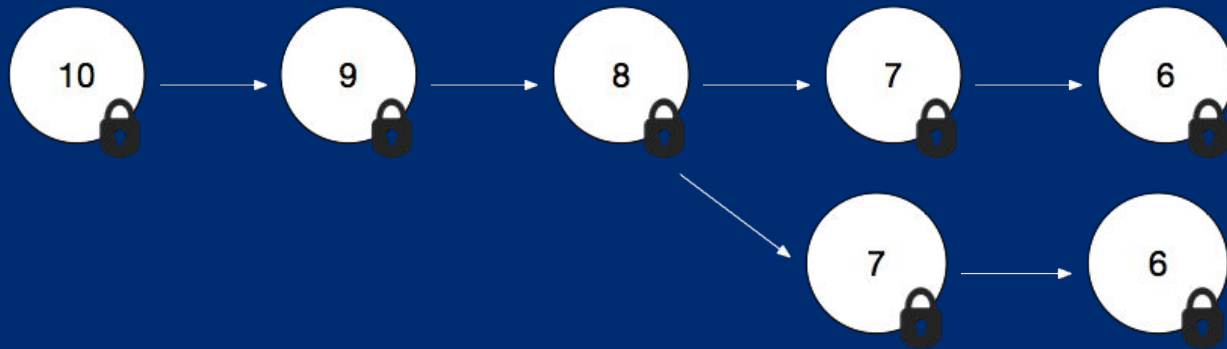
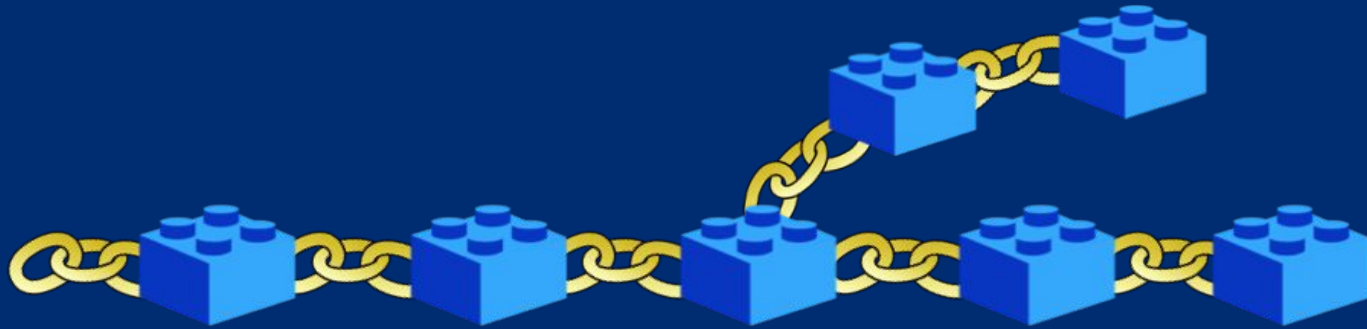


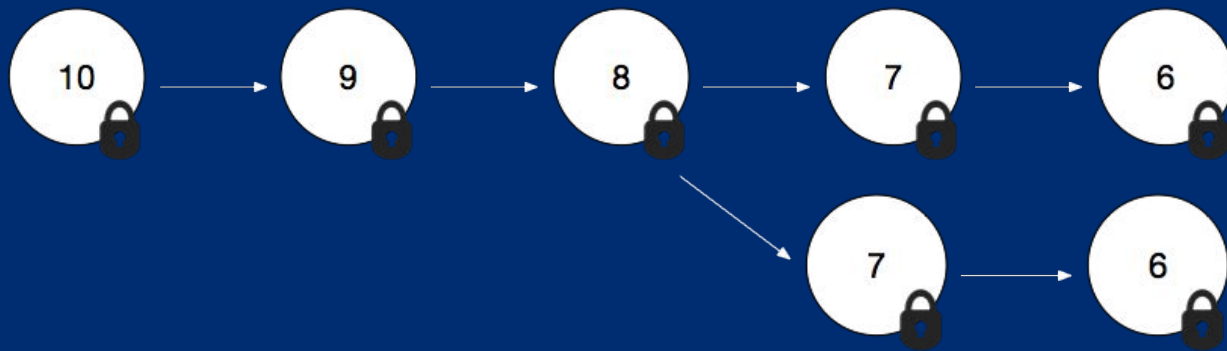
- Simplifying assumption: Single user ledgers

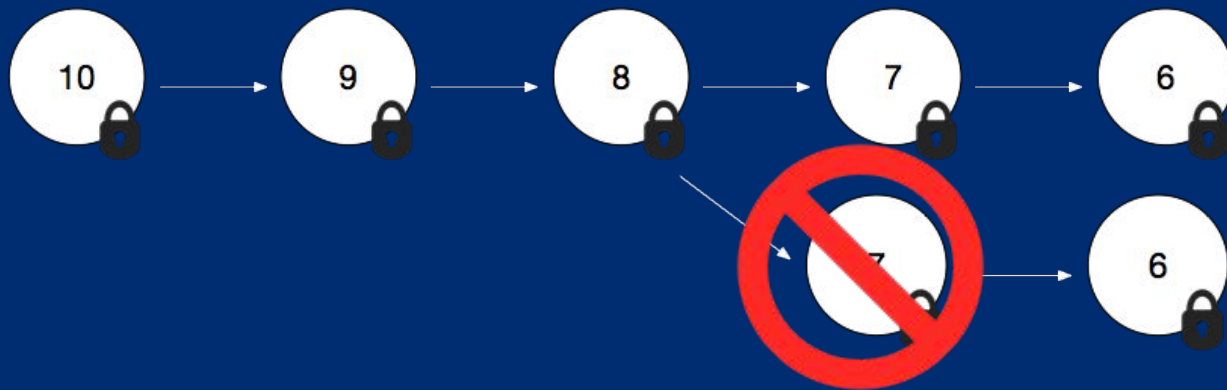








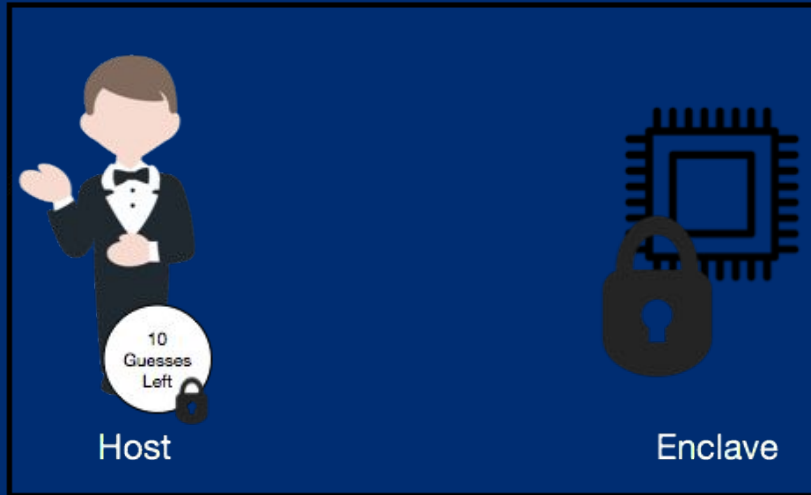




Ledger



User



Host

Enclave

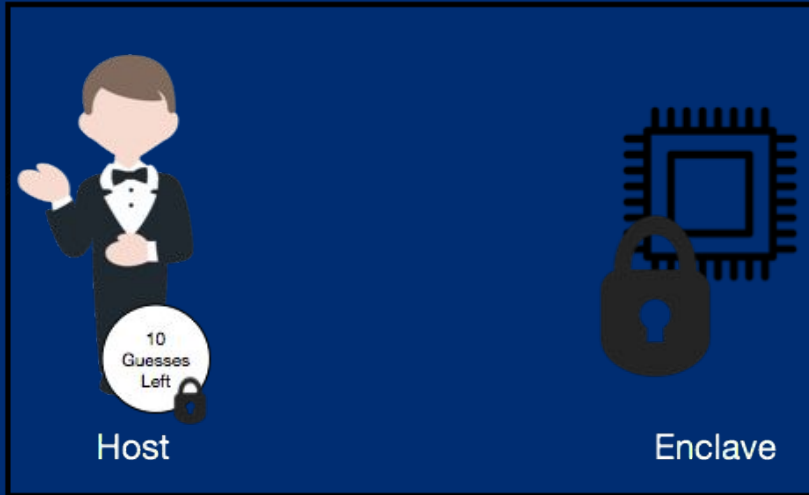


# Ledger



User

"1234"



Host

Enclave

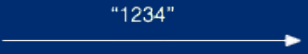
# Ledger



$$C = \text{Com}(10, "1234"; r)$$



User



Host

Enclave

# Ledger



Transaction Hash  
Previous Tx Hash  
C

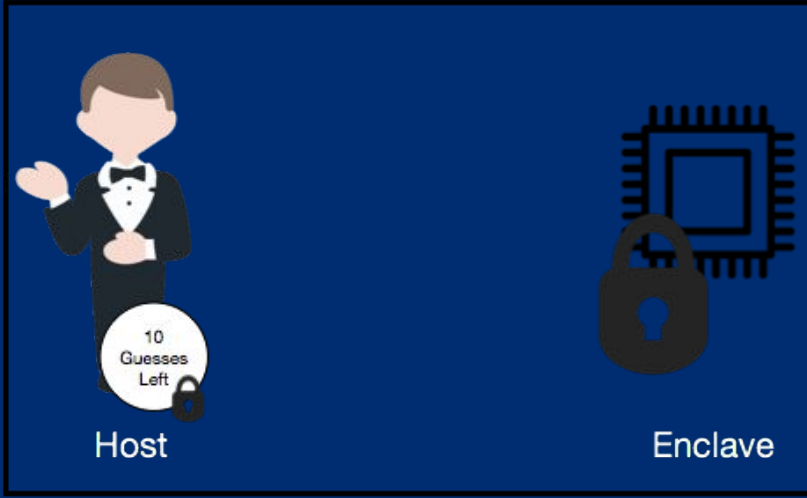


$$C = \text{Com}(10, "1234"; r)$$



User

"1234"



Host

Enclave

10  
Guesses  
Left

# Ledger



Transaction Hash  
Previous Tx Hash  
C

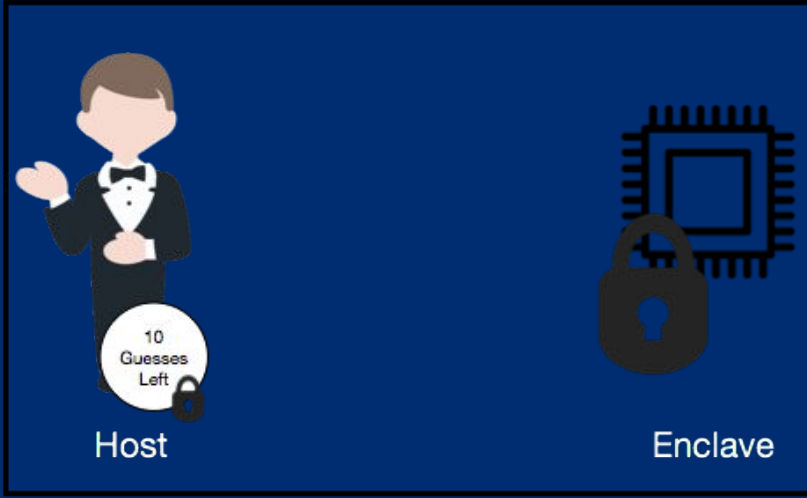


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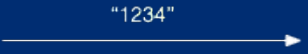
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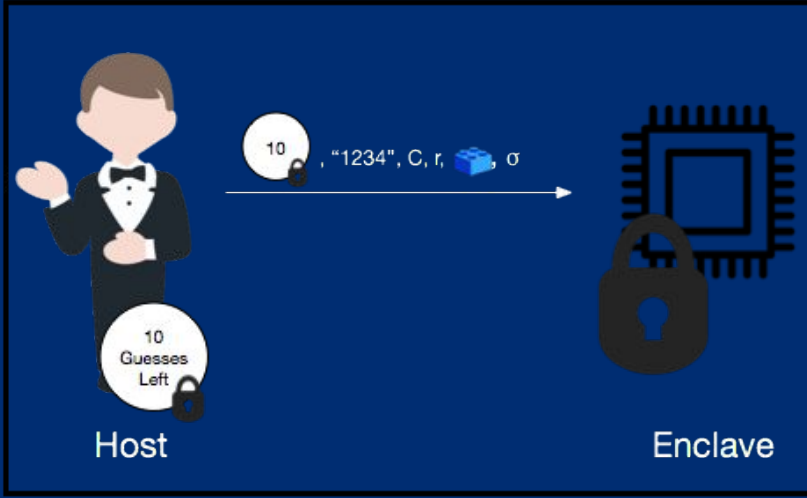
$\sigma$



User



"1234"



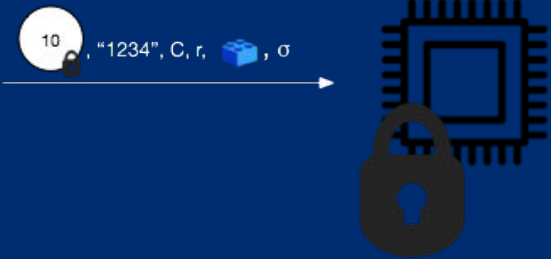
Host

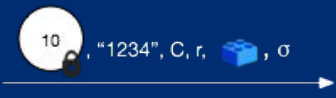
10

"1234", C, r, ,  $\sigma$



Enclave





Verify(,  $\sigma$ )

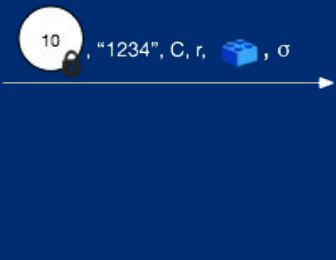
Verify( C = Com( , "1234"; r ) )



Derive encryption key for next state ←

Transaction Hash  
Previous Tx Hash  
C

Derive decryption key for previous state ←



Verify(,  $\sigma$ )

Verify(  $C = \text{Com}(10, "1234"; r)$  )

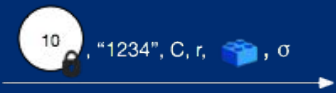




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C

Derive decryption key for previous state ←



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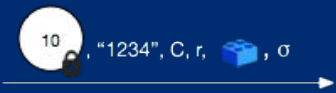
$10 = \text{Decrypt}( 10, \text{PRF}_{sk}(\text{Previous Tx Hash}) )$



Derive encryption key for next state ←

Transaction Hash  
Previous Tx Hash  
C

Derive decryption key for previous state ←



Verify(  ,  $\sigma$  )

Verify(  $C = \text{Com}( 10, "1234"; r )$  )

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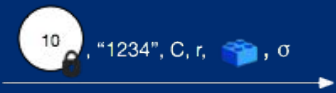
$9 \text{ key} = \text{Program}( 10, "1234"; \text{PRF}_{sk}(\text{"rand"} \parallel \text{Previous Tx Hash})$  )



Derive encryption key for next state ←

Transaction Hash  
Previous Tx Hash  
C

Derive decryption key for previous state ←



Verify(  ,  $\sigma$  )

Verify(  $C = \text{Com}( 10, "1234"; r )$  )

$10 = \text{Decrypt}( 10, \text{PRF}_{sk}(\text{Previous Tx Hash})$

$9 = \text{Program}( 10, "1234"; \text{PRF}_{sk}(\text{"rand"} \parallel \text{Previous Tx Hash})$

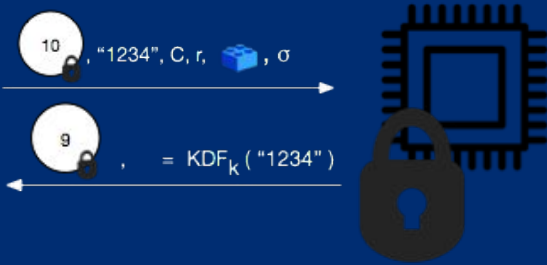
$9 = \text{Encrypt}( 9, \text{PRF}_{sk}(\text{Current Tx Hash})$



Derive encryption key for next state ←

Transaction Hash  
Previous Tx Hash  
C

Derive decryption key for previous state ←



Verify( [blue block] ,  $\sigma$  )

Verify(  $C = \text{Com}( 10 , \text{"1234"}; r )$  )

$10 = \text{Decrypt}( 10 , \text{PRF}_{sk} ( \text{Previous Tx Hash} )$

$9 = \text{Program}( 10 , \text{"1234"}; \text{PRF}_{sk} ( \text{"rand"} \parallel \text{Previous Tx Hash} )$

$9 = \text{Encrypt}( 9 , \text{PRF}_{sk} ( \text{Current Tx Hash} )$

# Ledger



Transaction Hash  
Previous Tx Hash  
C



$$C = \text{Com}(10, "1234"; r)$$

$\sigma$



User

"1234"



Host

10

9

10  
Guesses  
Left

"1234", C, r,  $\sigma$

$= \text{KDF}_K("1234")$



Enclave

# Ledger



Transaction Hash  
Previous Tx Hash  
C



$$C = \text{Com}(10, "1234"; r)$$

  $\sigma$

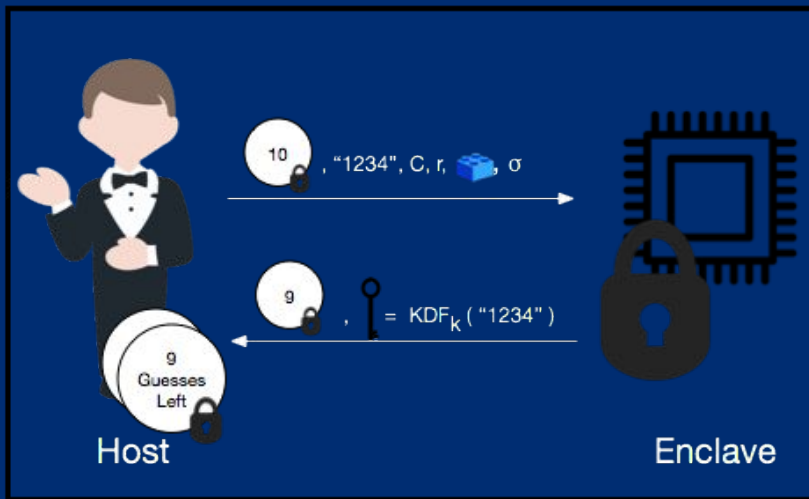


User

"1234"



Decryption Failure, 9 more



# Protocol Extensions

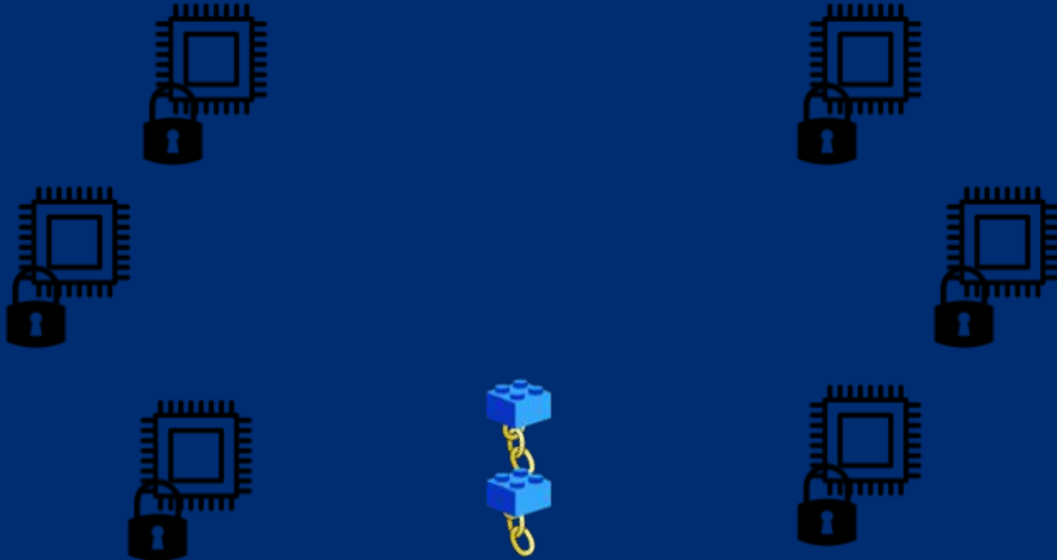
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- We have managed to condition execution on ledger postings
- Extension #1: Programs can require public posting
  - E.g. Error reporting, guaranteed logging
- Extension #2: One Time Programs
  - Swept under the rug: so far we have secure *multi-execution programs*
  - Derive *unique* valid hash chain from program code

# Additional Applications

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- Smart contracts computing on private data
  - Concurrent work with Intel's Private Data Objects
  - Later follow-up work in the same area [Eikiden]

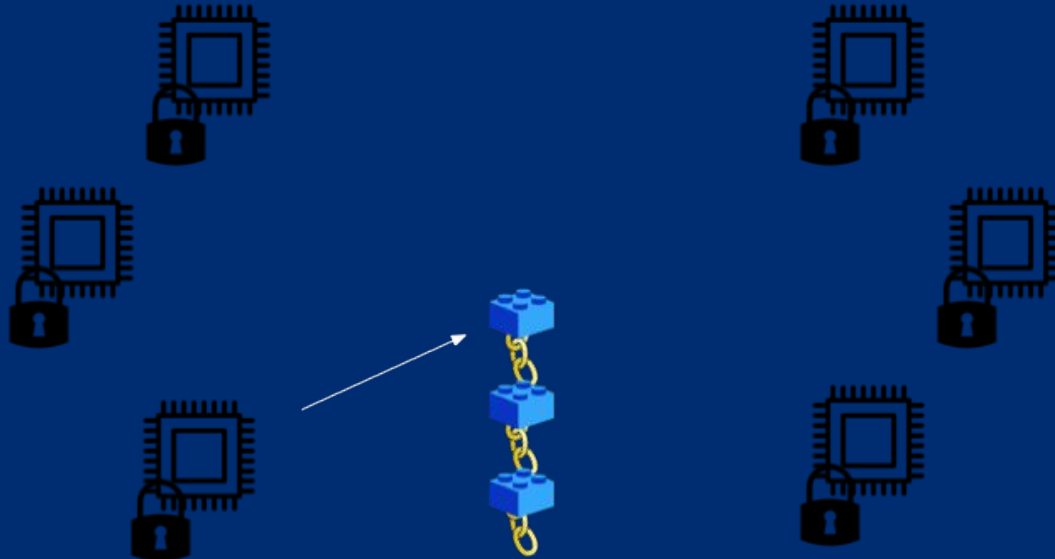




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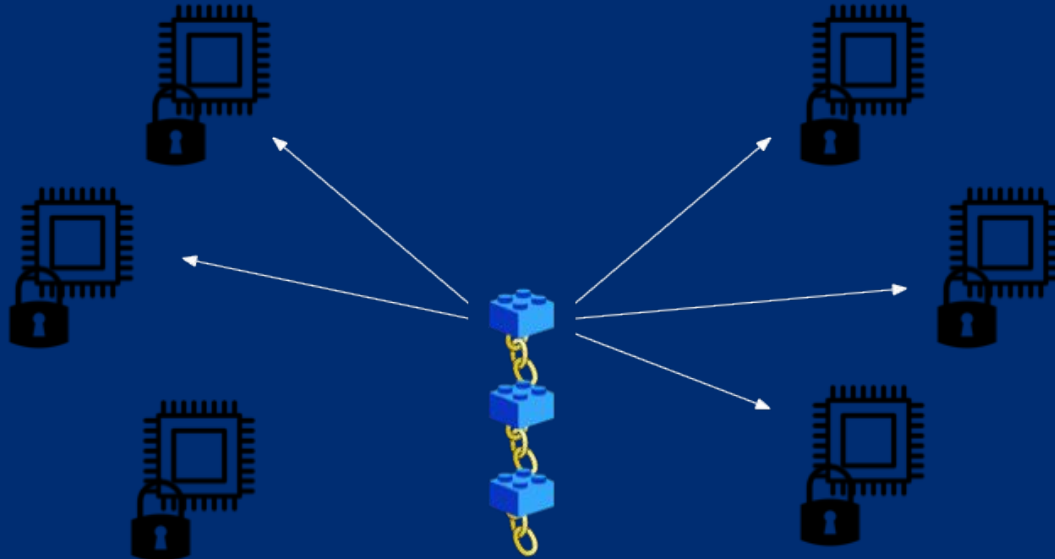
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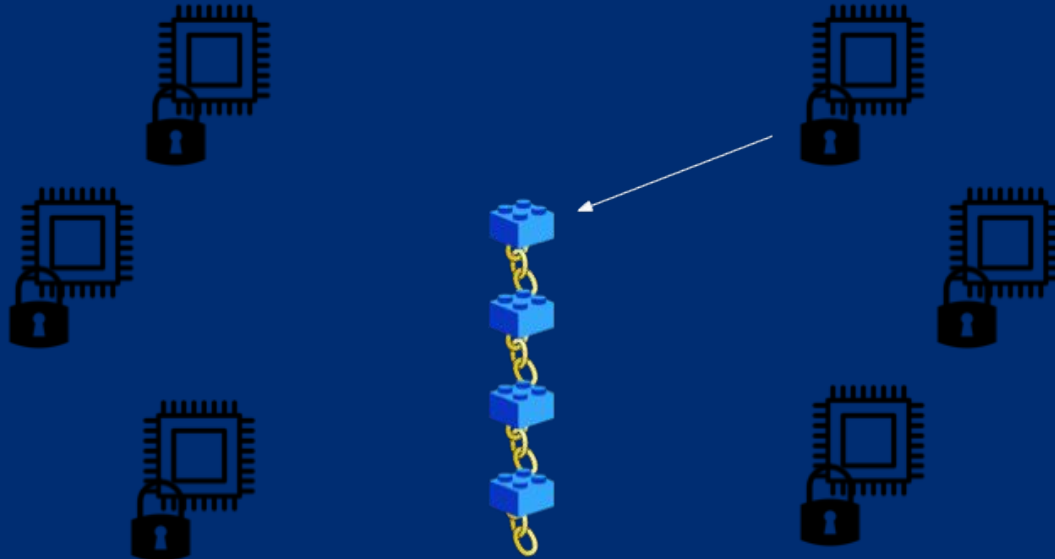
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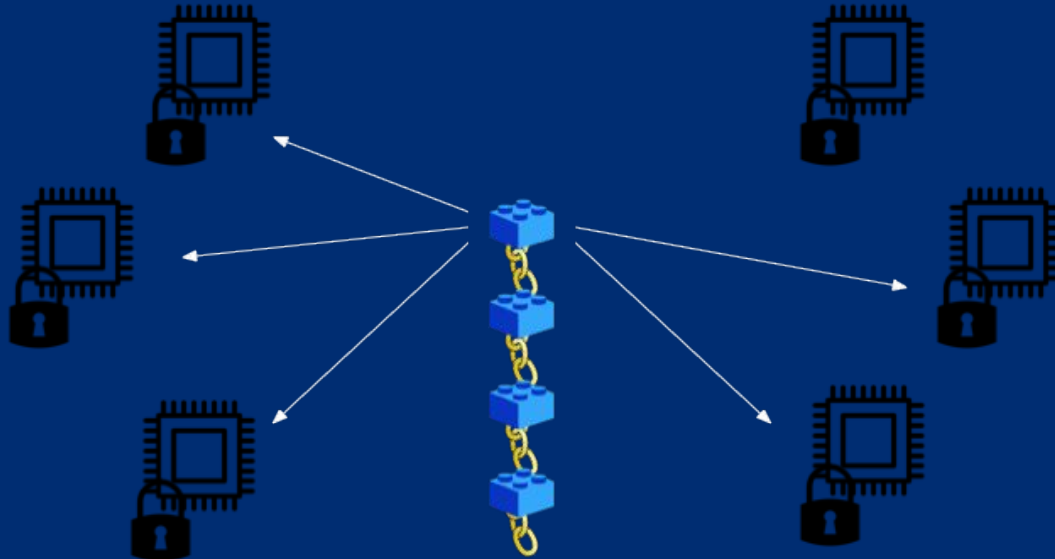
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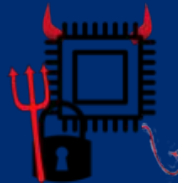


# Additional Applications

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- Autonomous Ransomware
  - Inevitable outcome of malicious trusted execution environments
  - Eliminates the need for command and control systems

Show me a valid  
cryptocurrency payment to my  
address and I'll give you the  
key!



# Conclusions

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- We create a novel protocol that provides trustworthy state for TEE's by binding state to an append-only ledger
- Ledgers are here to stay — lets do more than just currency-related research
- Keeping state is a difficult problem with wide ranging applications

# Thank You!

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gabe@kaptchuk.com

# Bonus Slides





