Automated Synthesis of Semantic Malware Signatures using MaxSAT



NDSS'17, San Diego

Newly Found Malware can steal bank details on Android phones



by **Ali Raza** 9 months ago



https://www.hackread.com/malware-can-steal-bank-details-android-phones/

Hundreds Of Operations Canceled After Malware Hacks Hospitals Systems

🛗 Thursday, November 03, 2016 🛛 🛔 Mohit Kumar



http://thehackernews.com/2016/11/hospital-cyber-attack-virus.html

Android Malware Used to Hack and Steal a Tesla Car By Catalin Cimpanu 📰 November 25, 2016 🛛 💓 06:05 AM 2

http://www.bleepingcomputer.com/news/security/android-malware-used-to-hack-and-steal-a-tesla-car/

37M

Total count of malware detected over 6 months

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Apposcopy Overview



Feng, et al. FSE'14

Apposcopy Overview

A high-level language for describing semantic properties of malware



Feng, et al. FSE'14

Apposcopy Overview

A high-level language for describing semantic properties of malware A novel static analysis for deciding if an app matches the signature of a family



Feng, et al. FSE'14



Caveats



Writing signatures is tedious

Caveats

Monkey Test & Time Service



Writing signatures is tedious



Vulnerable to semantic obfuscation

Goal



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• Infer a signature from few samples of a malware family



Goal

- Infer a signature from few samples of a malware family
- Approximate matching algorithm that is resistant to semantic obfuscation



Inter-Component Call Graph



Inter-Component Call Graph





Activity1

CIOSCUD

Android

Activity2 Inter-Component Call Graph

 Image: Stratume
 Image: Stratume

 Image: Stratume
 Image: Stratume



CIOECUD

Android













Solid line: control property Dashed line: data property



GDEvent(SMS_RECEIEVED). GDEvent(NEW_OUTGOING_CALL). GoldDream :- receiver(r), icc(SYSTEM, r, e, _), GDEvent(e), service(s), icc*(r, s), flow(s, DeviceId, s, Internet), flow(s, SubscriberId, s, Internet).

GoldDream Signature



Service1 **DeviceId -> Internet** Activity1 Activity2 Inter-Component Call Graph nten CIOSCUD ContentProvider **Receiver1** Android

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Component

Predicate

GoldDream Signature

Solid line: control property Dashed line: data property



Component
PredicateControl
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GoldDream Signature







Given *n* malware samples from family *F*, compute its signature *S*



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Any signature that matches n samples



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Empty signature could also be a solution!


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10







10















<u>MaxSat</u>: Given a UNSAT boolean formula in CNF, determine the <u>maximum</u> number of satisfied clauses $(x_0 \lor x_1) \land (\neg x_0 \lor x_1) \land (x_0 \lor \neg x_1) \land (\neg x_0 \lor \neg x_1)$

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<u>Soft Clause</u>: preferable to be satisfied but could be UNSAT. Each has different <u>weight</u> since some are more important than the others

Find an assignment s.t. the total weight of satisfied clauses is maximized

 $\{x_0 \mapsto 0, x_1 \mapsto 0\}$

• <u>Hard Clause</u>: common subgraph (control-flow property)

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- <u>Weight</u> for each clause
 - Inverse frequency from benign samples
 - Higher weight to features that are commonly found in malware

$$\mathcal{O} = \sum_{v,v' \in V} x_0(v,v') + \sum_{v,v' \in V} \sum_{d \in \mathcal{D}} w_{(v,v',d)} y_0(v,v',d).$$

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Control properties

 $\int w_{(v,v',d)} y_0(v,v',d).$ $x_0(v,v') +$ $\mathcal{O} =$ Hard $v, v' \!\in\! V \; d \!\in\! \mathcal{D}$ $v, v' \in V$



















Approximate matching

Now that we have the signature...

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Utilize existing signature inference algorithm to decide if a sample A belongs to a family F:
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$$\delta(\mathcal{A}, \mathcal{F}) = \frac{f(\text{INFERSIGNATURE}(\mathcal{A}, \mathcal{S}_{\mathcal{F}}))}{f(\mathcal{S}_{\mathcal{F}})}$$

f(S): Weighted sum of the number of nodes and edges in S











Resistant to semantic obfuscation!

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- RQ2: How effective is Astroid at detecting zeroday malware?
- RQ3: How does Astroid compare against stateof-the-art malware detectors?













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Our approximate matching is effective!









False positive rate: Drebin(NDSS'14): 1%, MassVet (Security'15): 175/503, Astroid: 0.04%

Astroid achieves high detection rate with low FP!



Conclusion



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 Automatically infer semantic malware signature from very few samples



Conclusion

- Automatically infer semantic malware signature from very few samples
- Our approximate matching is resilient to semantic obfuscations



Thank you!

Automated Synthesis of Semantic Malware Signatures using Maximum Satisfiability. <u>Yu Feng</u>, Osbert Bastani, Ruben Martins, Isil Dillig, Saswat Anand. NDSS 2017.

EXPLORER: Query- and Demand-Driven Exploration of Interprocedural Control Flow Properties. <u>Yu Feng</u>, Xinyu Wang, Isil Dillig, Calvin Lin. OOPSLA 2015.

Apposcopy: Semantics-Based Detection of Android Malware through Static Analysis. <u>Yu Feng</u>, Saswat Anand, Isil Dillig, Alex Aiken. FSE 2014.