
Address-Oblivious Code Reuse: On the Effectiveness of Leakage-Resilient Diversity

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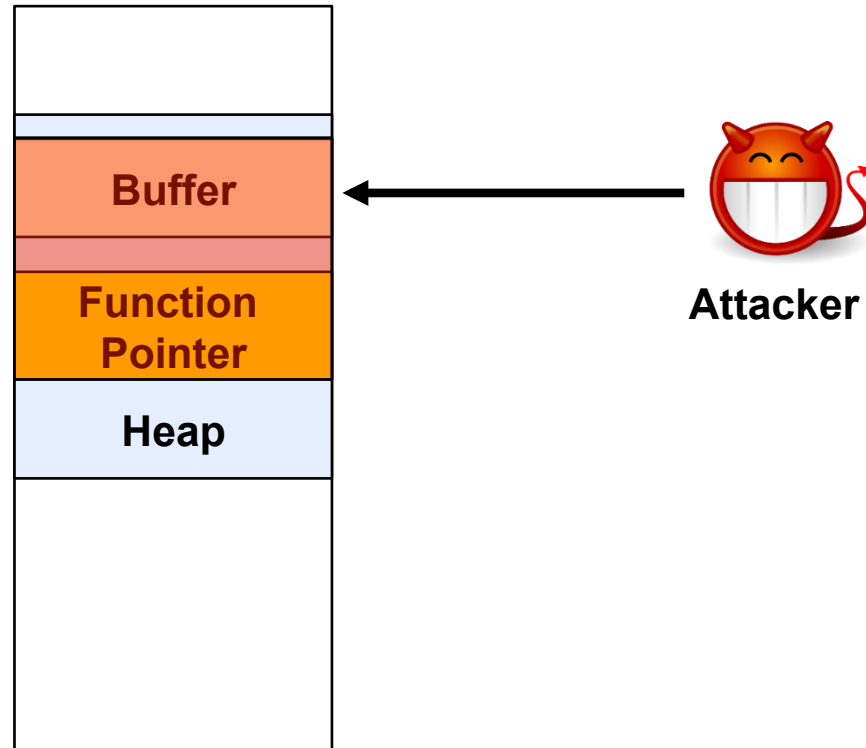


Bottom-Line Upfront

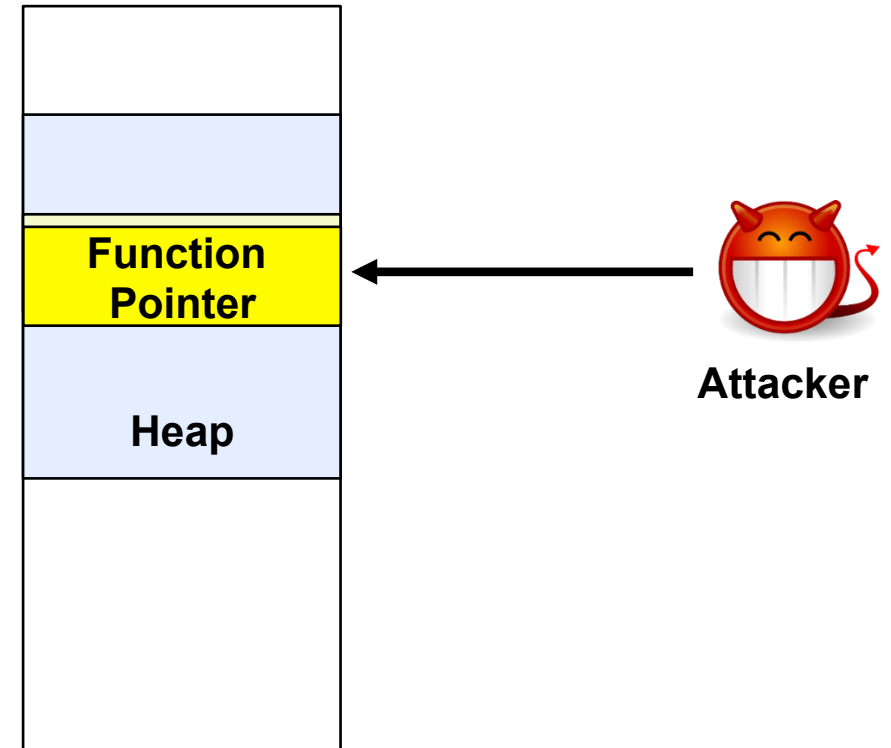
- Code diversity techniques are vulnerable to information leakage
- Recent leakage-resilient techniques employ “execute-only” memory permissions to prevent information leakage
- We present a generic type of attack called Address-Oblivious Code Reuse (AOCR) that can bypass recent leakage-resilient techniques
- We provide 3 real-world exploits



Memory Corruption Attacks



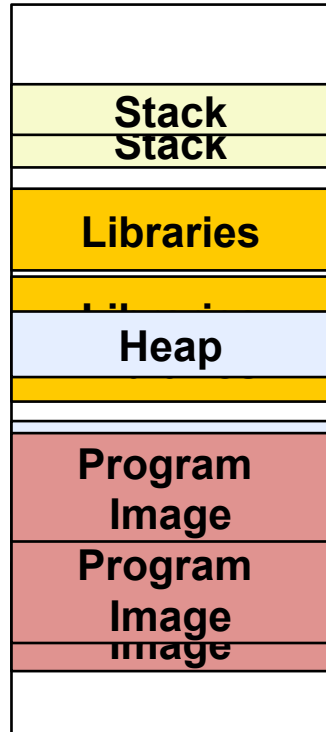
Spatial Memory Violation



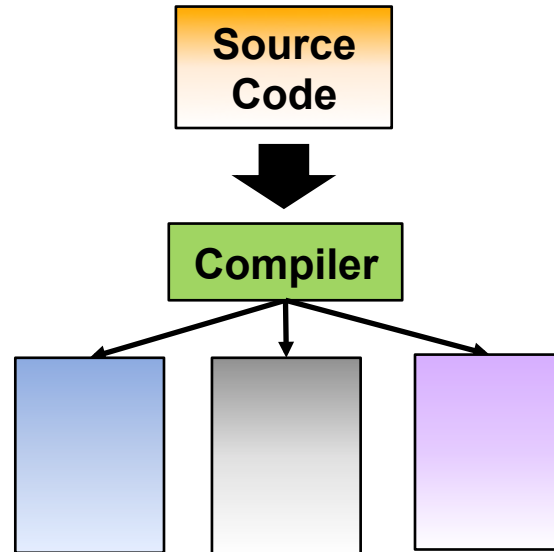
Temporal Memory Violation



Code Diversification Techniques

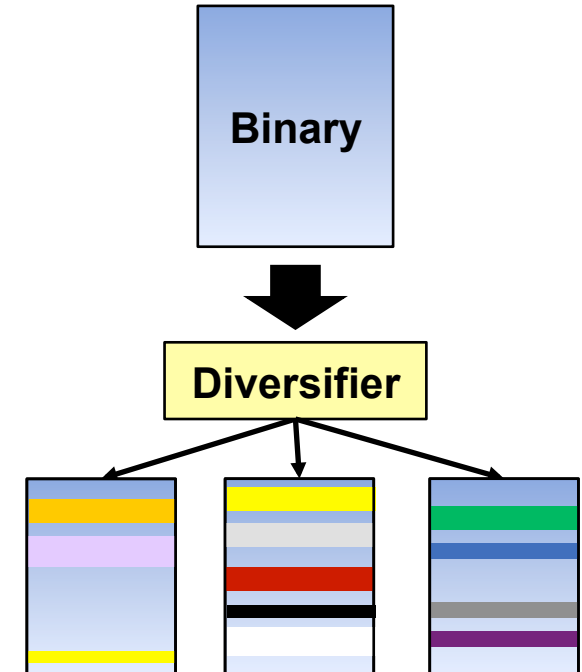


Address Space Layout Randomization (ASLR)



Diverse binaries

Compile-Time Diversity †

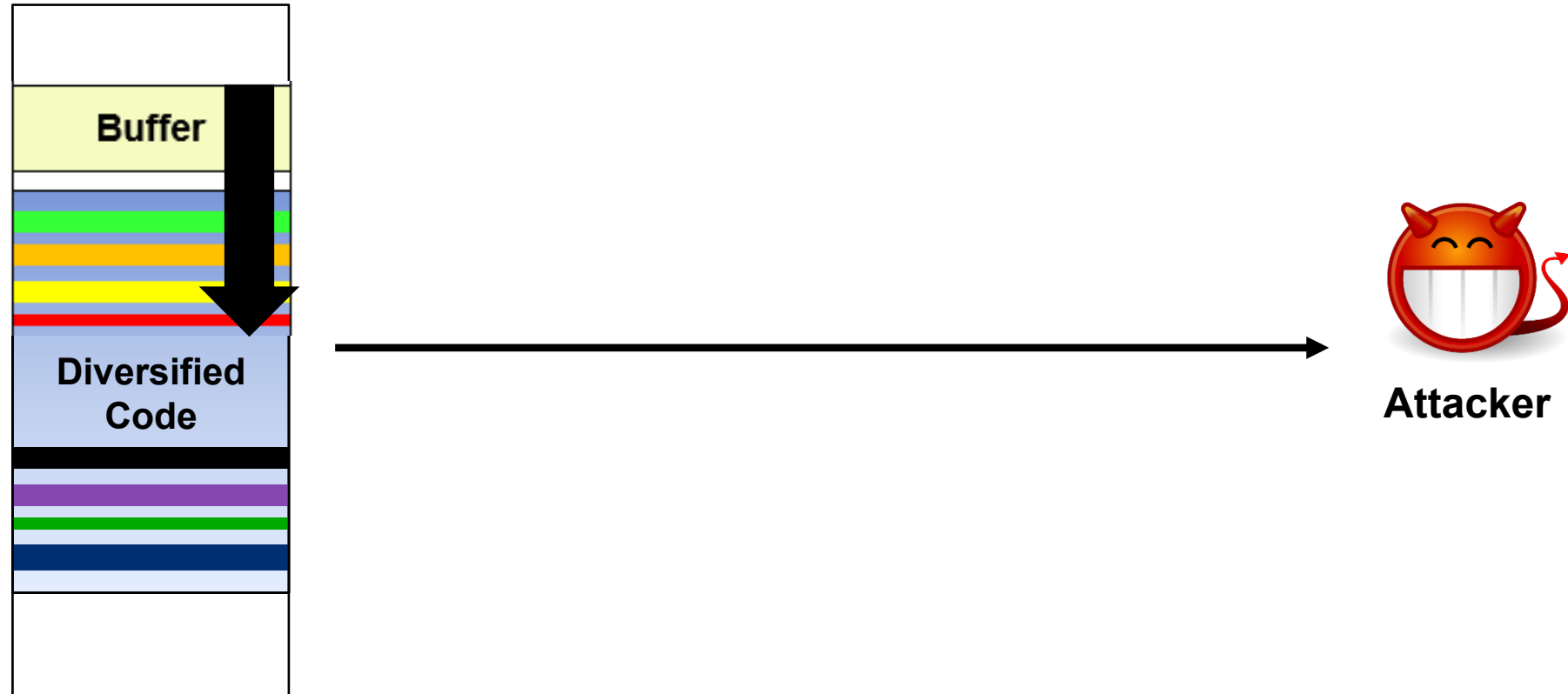


Diverse binaries

Binary Rewriting ‡

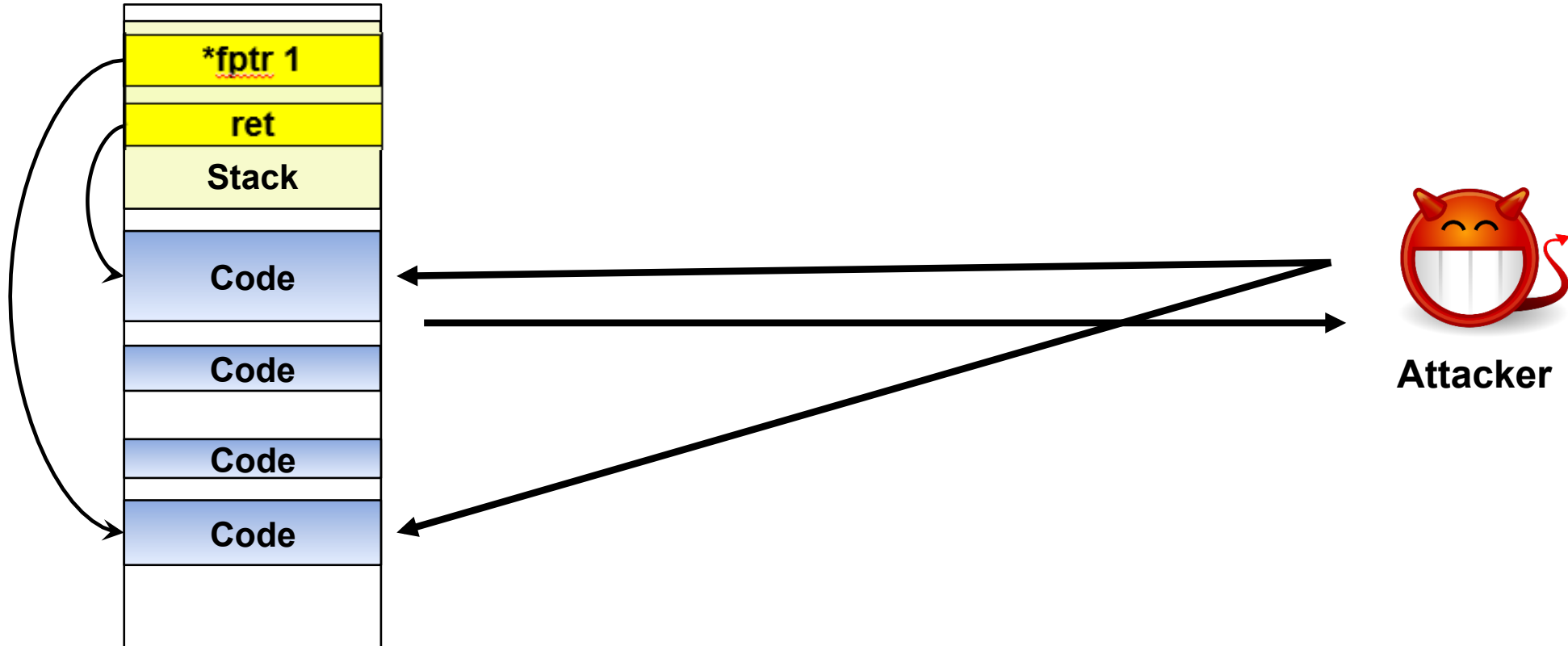


Information Leakage (Direct Memory Disclosure)



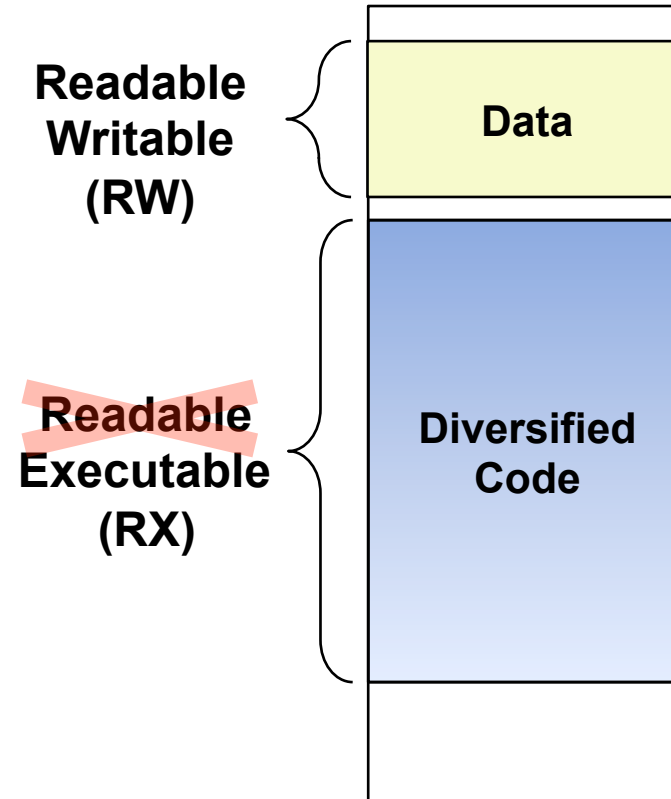


Information Leakage (Indirect Memory Disclosure)



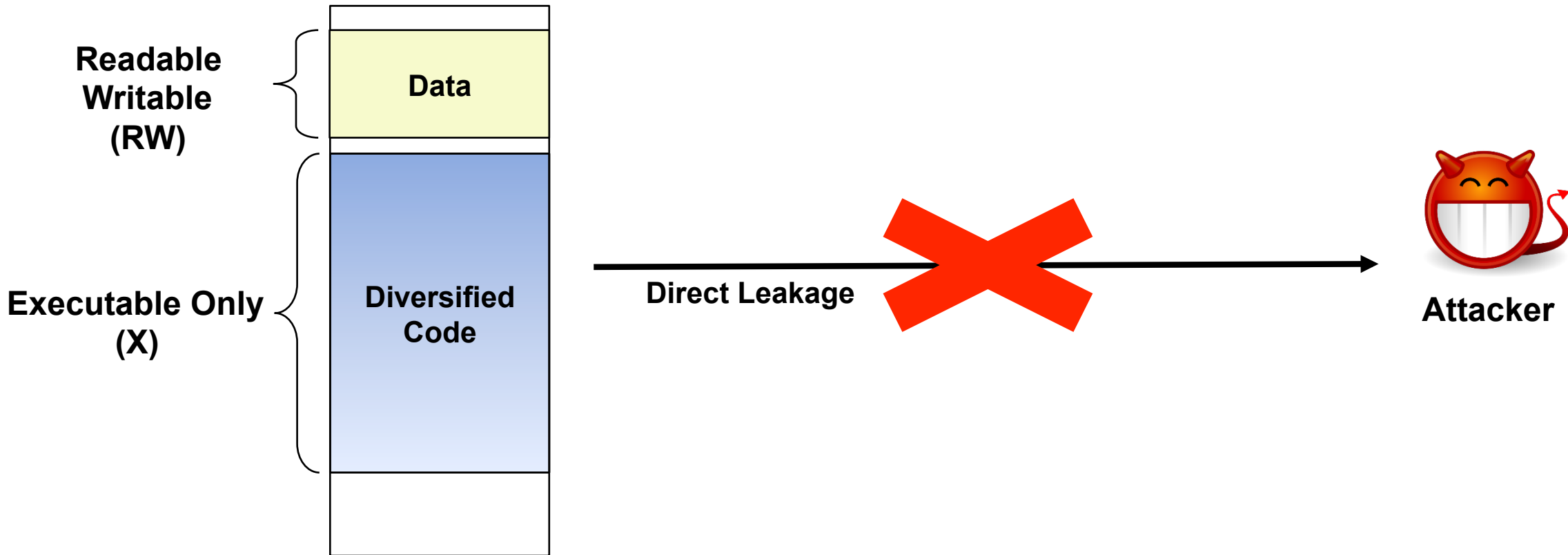


Memory Permissions



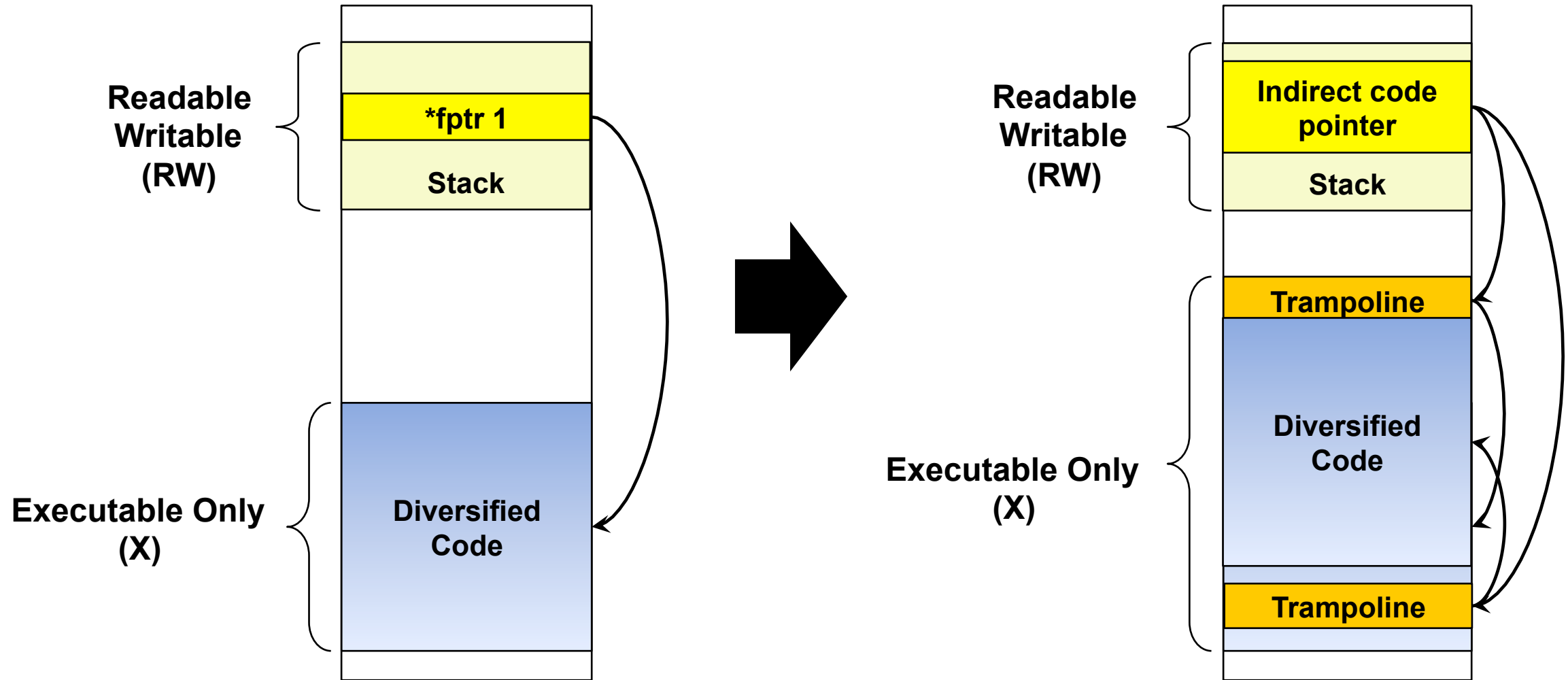


Leakage Resilient Diversity





Indirect Leakage Prevention





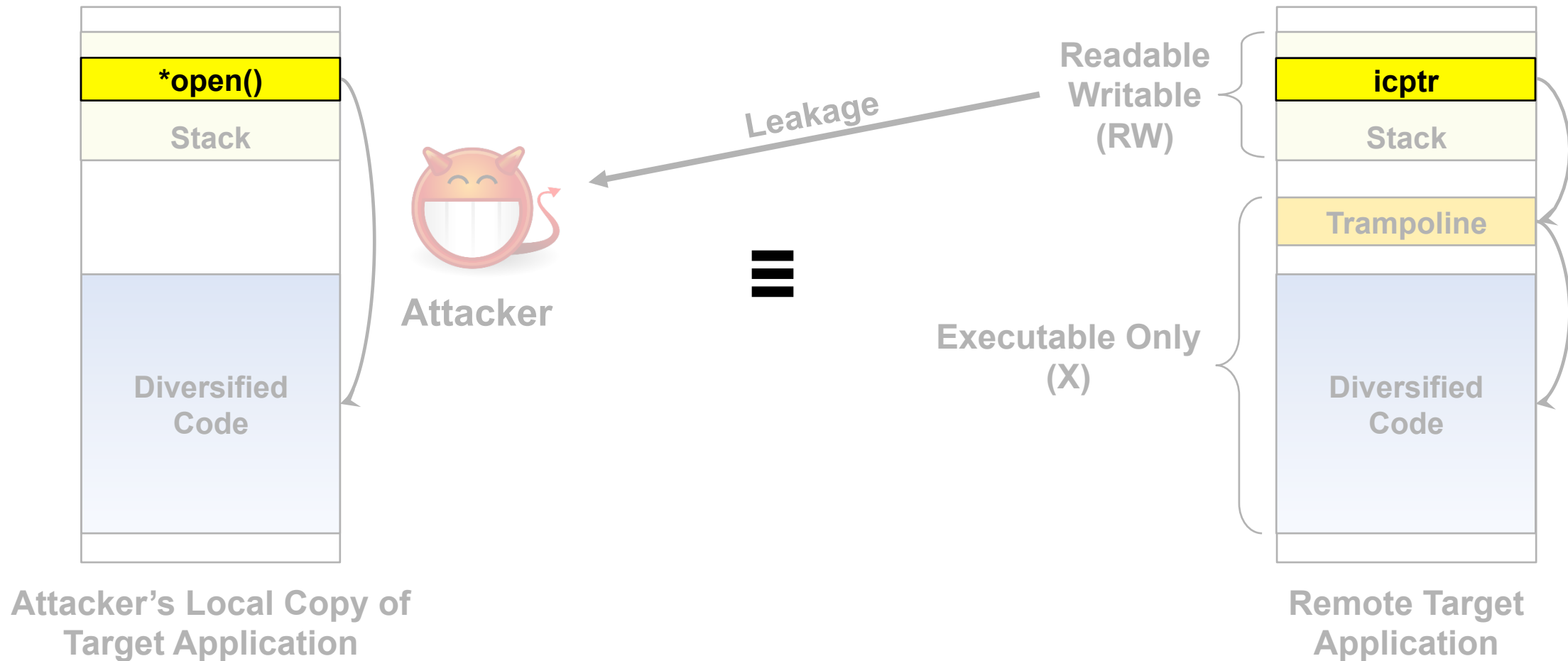
Research Questions

- **Indirect code pointers create a surrogate for code**
- **Can attackers reuse code at the granularity of indirect code pointers?**
- **Can they accurately identify the corresponding functions?**
- **Can they chain indirect code pointers together?**



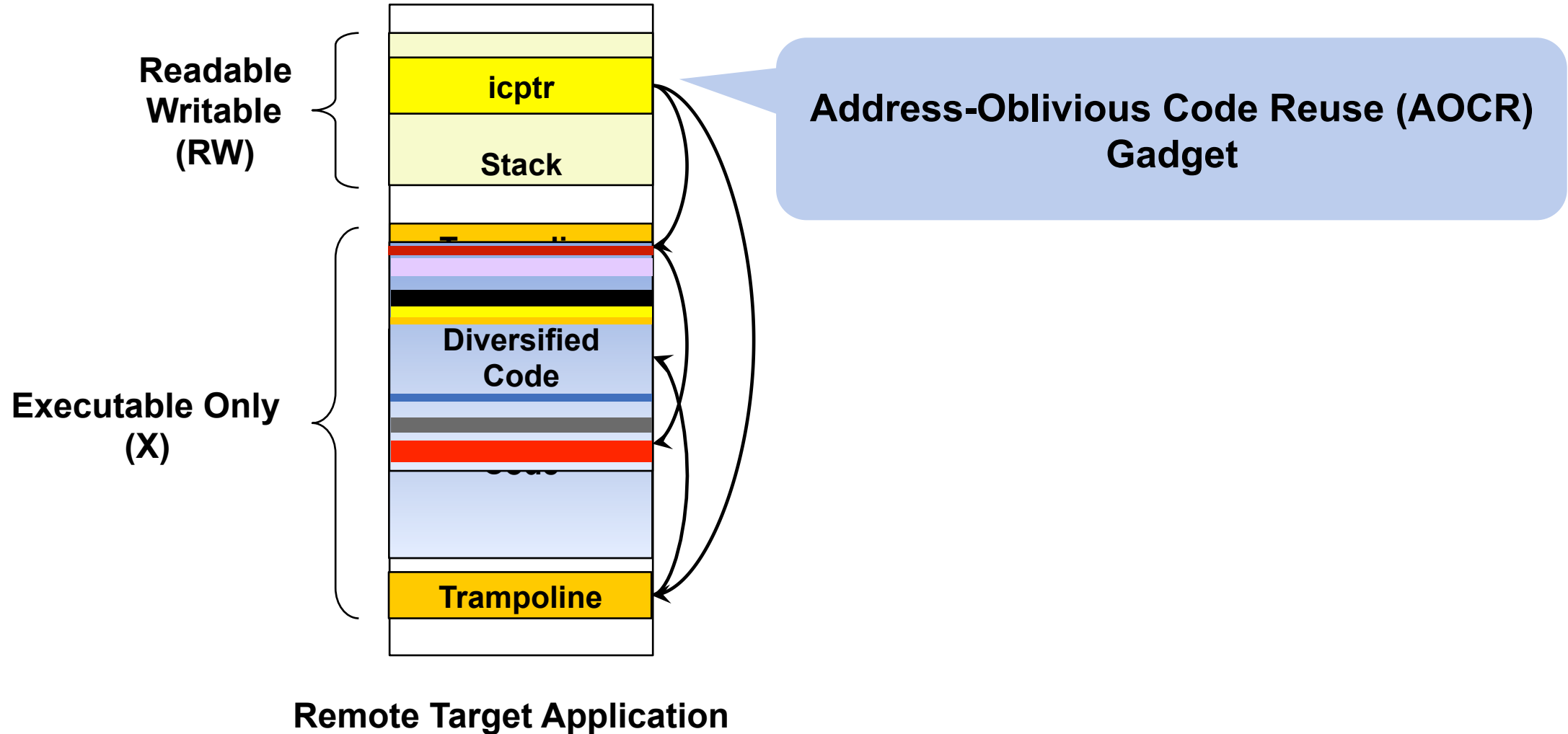
Profiling Indirect Code Pointers

- Goal: identify the function corresponding to each indirect code pointer



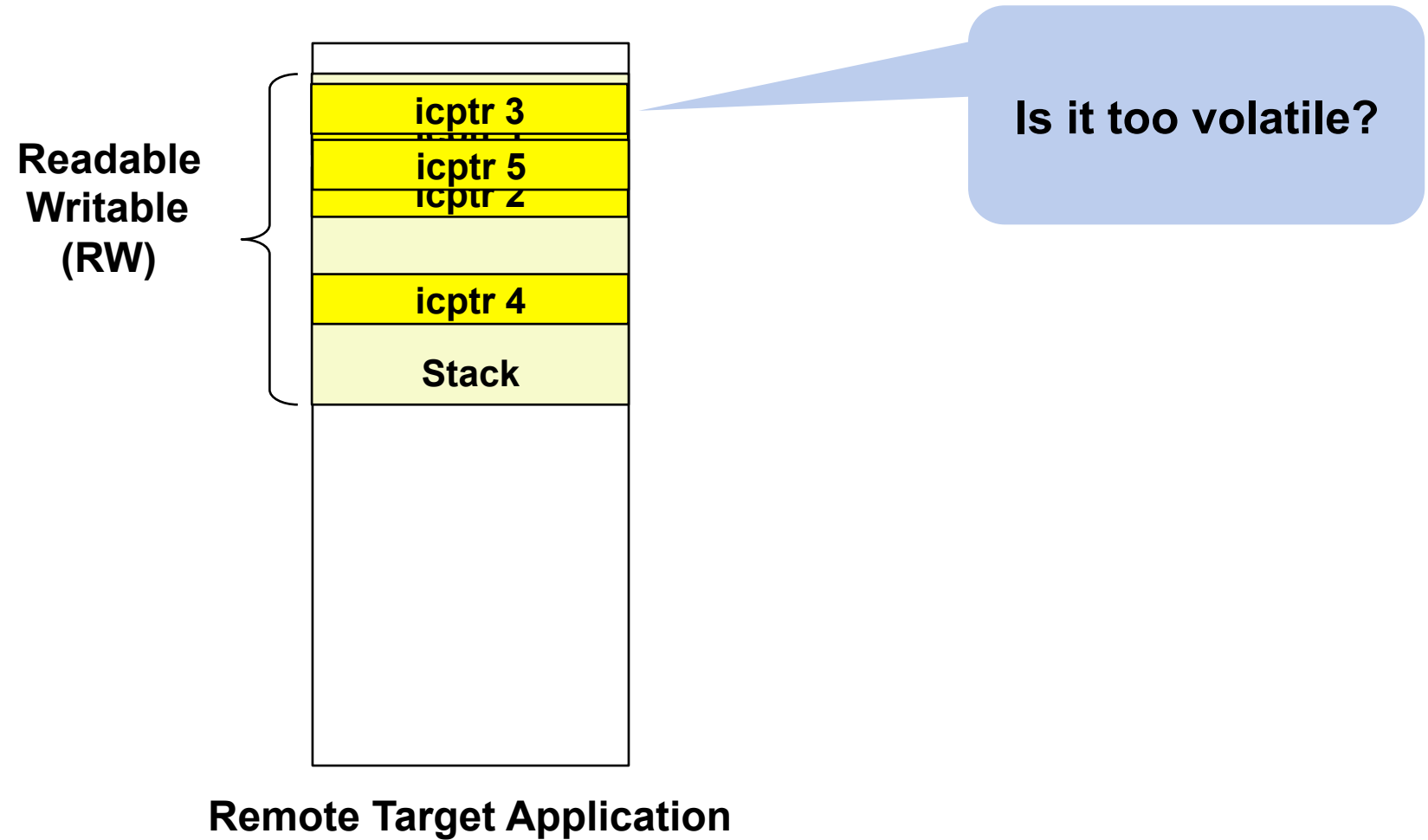


Address-Oblivious Code Reuse





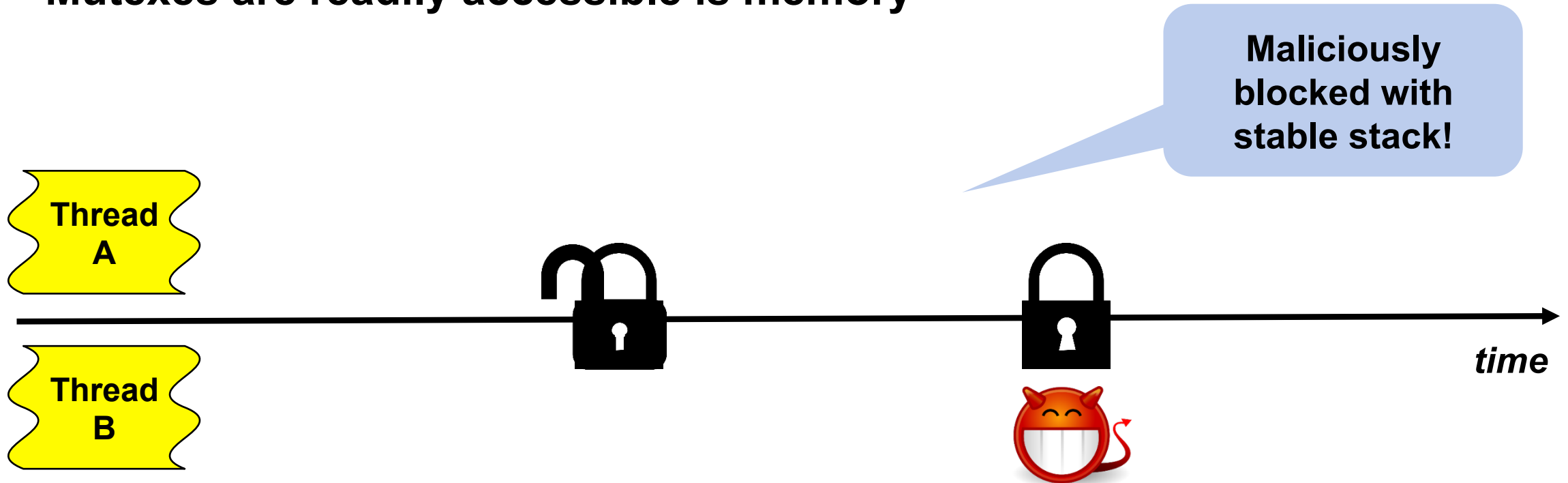
Accurate Profiling





Accurate Profiling using Malicious Thread Blocking (MTB)

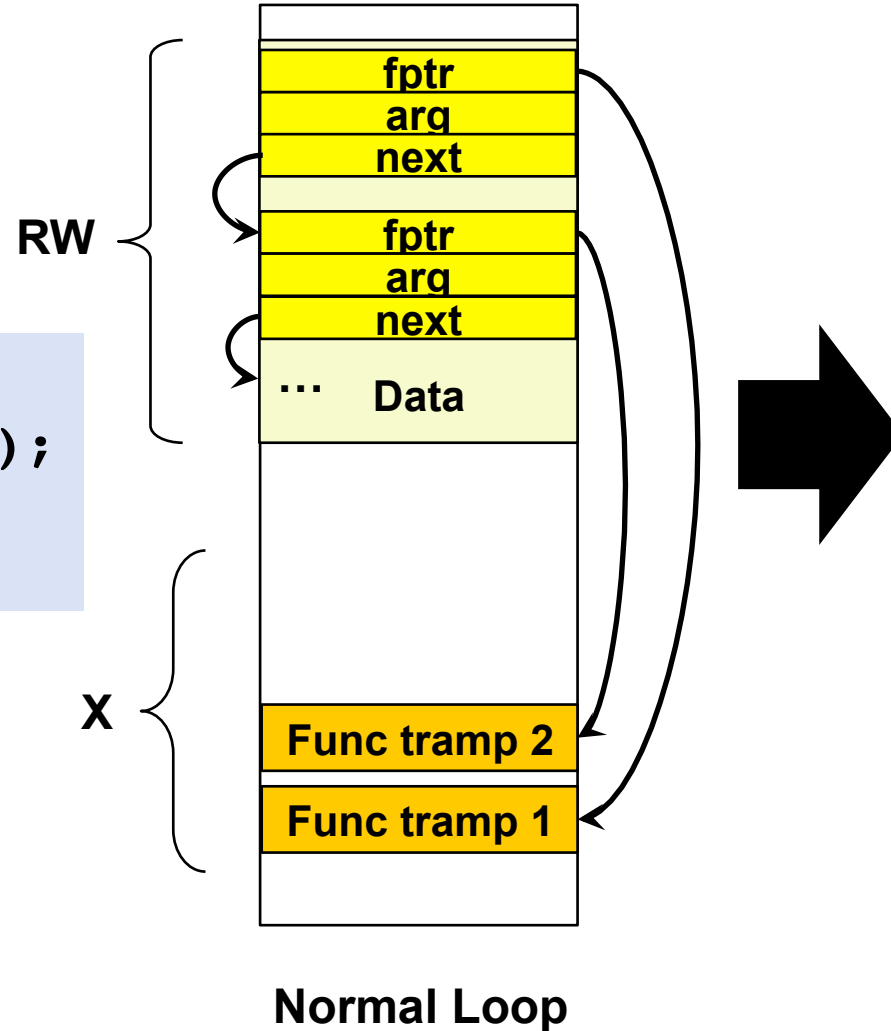
- A thread can force another thread to halt by maliciously setting a mutex
- Mutexes are readily accessible in memory





Chaining Gadgets Together using Malicious Loop Redirection (MLR)

```
while (task) {  
    task->fptr(task->arg);  
    task = task->next;  
}
```





Nginx Proof-of-Concept Exploit

- 1. Locate a mutex for MTB**
- 2. Profile an indirect code pointer for open (1st AOCR gadget)**
- 3. Profile an indirect code pointer for `_IO_new_file_overflow` (2nd AOCR gadget)**
- 4. Corrupt Nginx's task queue to call our profiled trampolines using MLR**



Implementation Challenges of Execute-Only Permissions

- **Forged Direct Memory Access (FDMA)**
 - A malicious application forges a software-based DMA call to kernel
 - Uses `O_DIRECT` flag in Linux
 - DMA request bypasses memory permissions

- **Procfs**
 - Ubiquitous facility in Linux
 - Provides memory maps and addresses
 - Blocking it breaks many benign applications
 - Protections such as GRSecurity's permissions will not block it



Impact on Leakage-Resilient Diversity Techniques

	Direct Leak		Indirect Leak	
	TLB-mediated (Buffer Overread)	Non-TLB-mediated (Forged DMA)	Code Pointer Leak (Ret address leak)	Indirect Code Pointer Leak (AOCR)
PointGuard				
Oxymoron				
Isomeron				
XnR				
HideM				
Readactor				
Heisenbyte				
NEAR				
ASLR-Guard				
TASR				



Possible Countermeasures

- **Complete memory safety**
- **Data randomization**
- **Authentication of indirect calls and returns**
 - **Use HMAC tokens to disallow redirection of indirect code pointers**
 - **Similar to cryptographically-enforced CFI (CCFI)**



Conclusion

- **Code pointers pose a major challenge to leakage-resilient diversity**
- **AOCR attacks bypass code pointer obfuscation by profiling indirect code pointers**
- **Malicious threat blocking (MTB) allows accurate profiling**
- **Malicious loop redirection (MLR) allows chaining AOCR gadgets**
- **Effective defenses should incorporate aspects of diversification and enforcement**



Questions?