

# On Limitations of Designing Leakage-Resilient Password Systems: Attacks, Principles and Usability

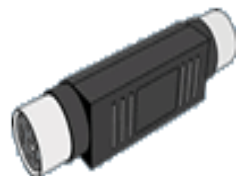
**Qiang Yan**, Jin Han, Yingjiu Li, Robert H. Deng

*School of Information Systems*

*Singapore Management University*

# Leakage-Resilient Password Systems (LRPS)

- Malware, e. g. **software keylogger**, **MITM-at-the-browser**
- Untrusted input device  
e.g. **hardware keylogger**



- Shoulder surfing  
e. g. **hidden camera recording**



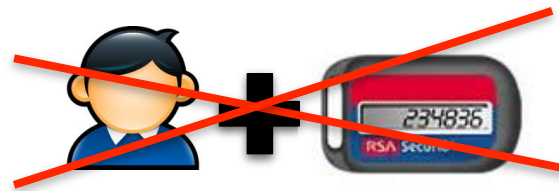
# Leakage-Resilient Password Systems (LRPS)

- Assumption
  - strong passive attacker



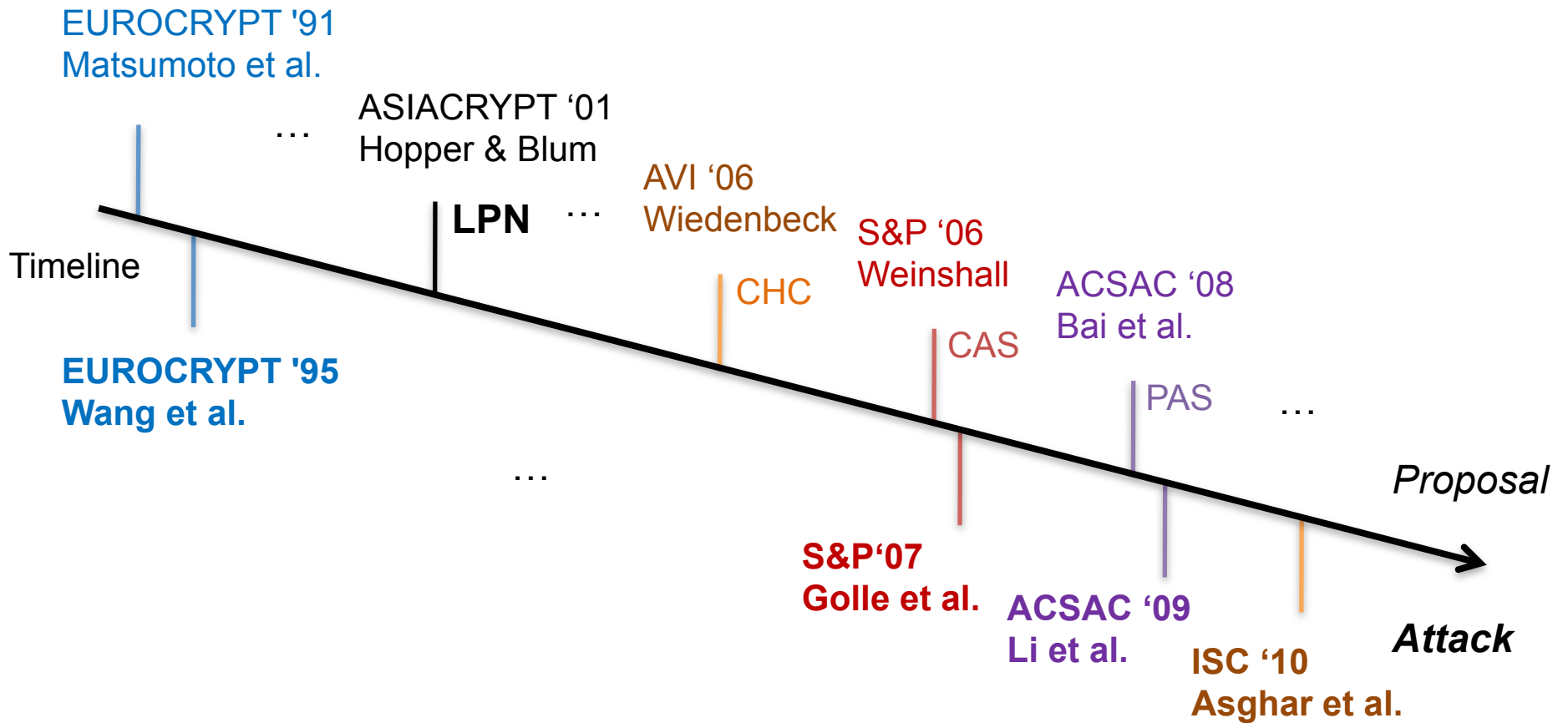
*Adversary sees **everything** below the red line.*

- Unaided user



Aided User

# Prior efforts on LRPS for unaided humans



# The $k$ -out-of- $n$ LRPS Paradigm

User's **root secret** (i.e. password) consists of  $k$  secret elements out of  $n$ .



User

1. **Challenge** with a window size  $w$  generated based on a **round secret** (i.e. a portion of root secret)



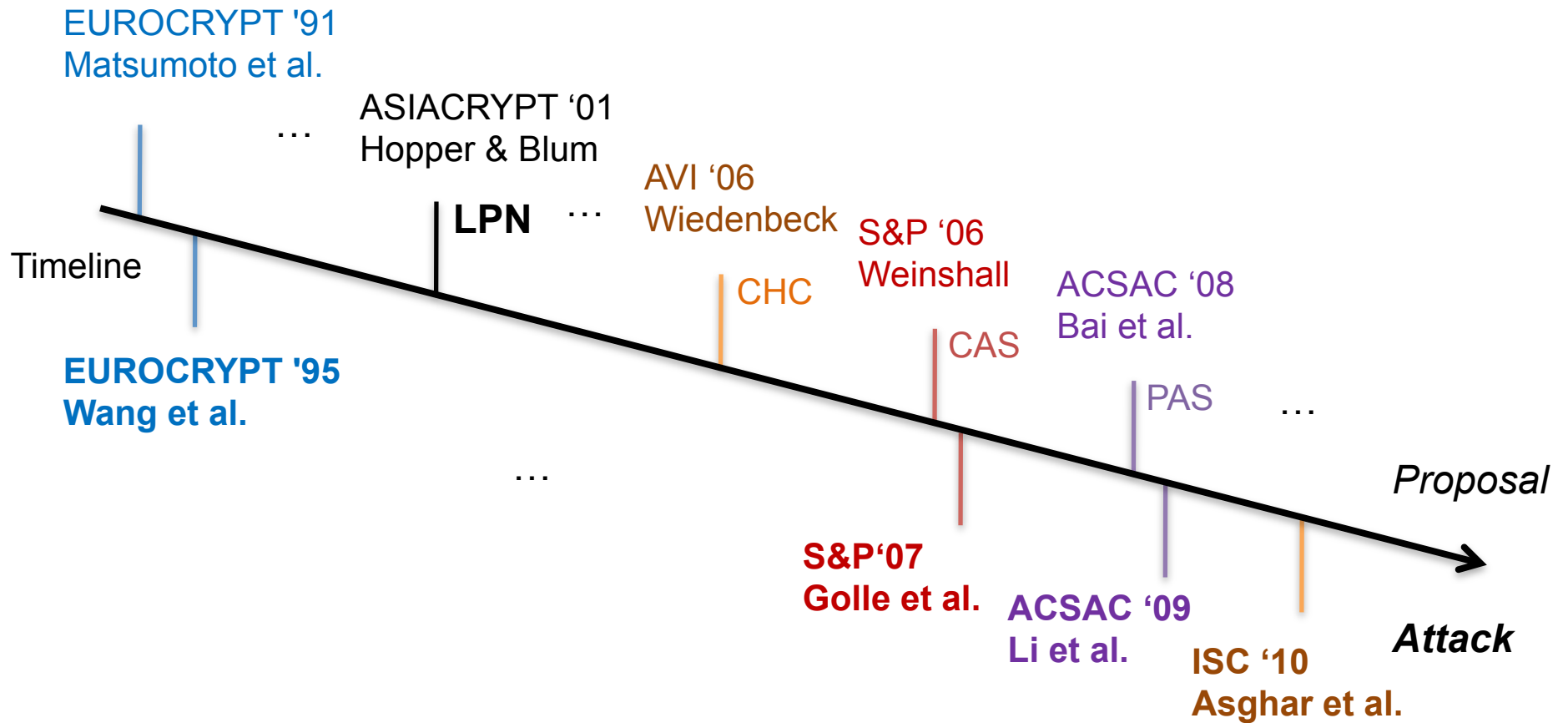
2. **Response** based on the knowledge of root secret



Server

*Repeat steps 1 and 2, until the number of correct user responses reaches a **threshold**.*

# Prior efforts on LRPS for unaided humans



# Two generic attacks

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- **Brute force**
  - Eliminate password candidates that **do not lead to correct responses**.
  - Effectiveness is *design-independent*.
    - Applicable to any LRPS with small password space
- **Statistical analysis**
  - Find out the **most likely** passwords.
  - Effectiveness is *design-dependent*.
    - Applicable to many LRPSs even with large password space
- They are common knowledge but are **underestimated**.

# Statistical bias in decision paths (1/2)

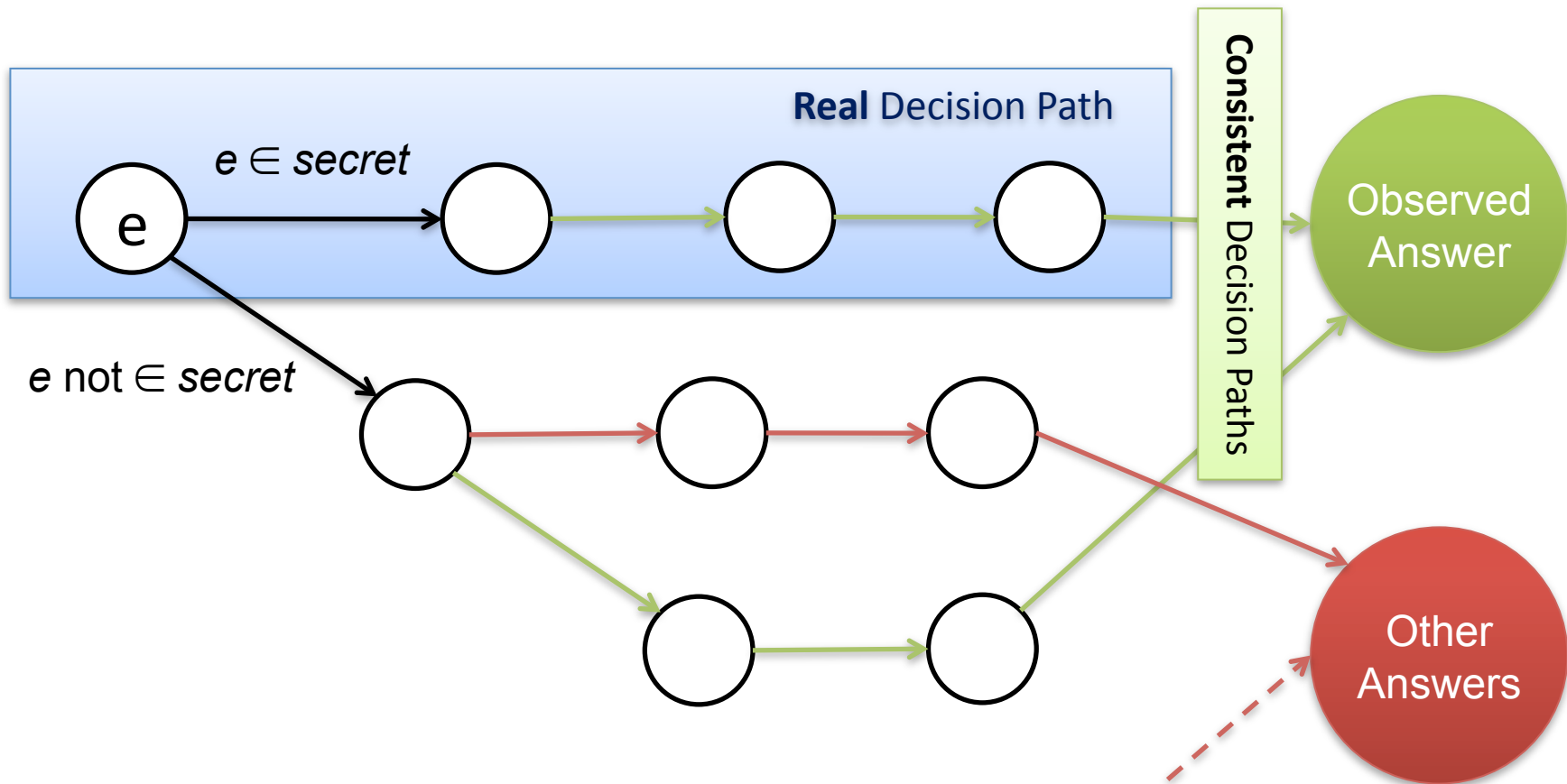
CAS High [S&P06, Weinshall]: Root secret consists of  $k = 30$  images out of 80



1. Start from upper-left corner.
2. Move down if the current image is a secret image; Otherwise move right.
3. Answer = the number associated with the exit.



# Probabilistic decision tree

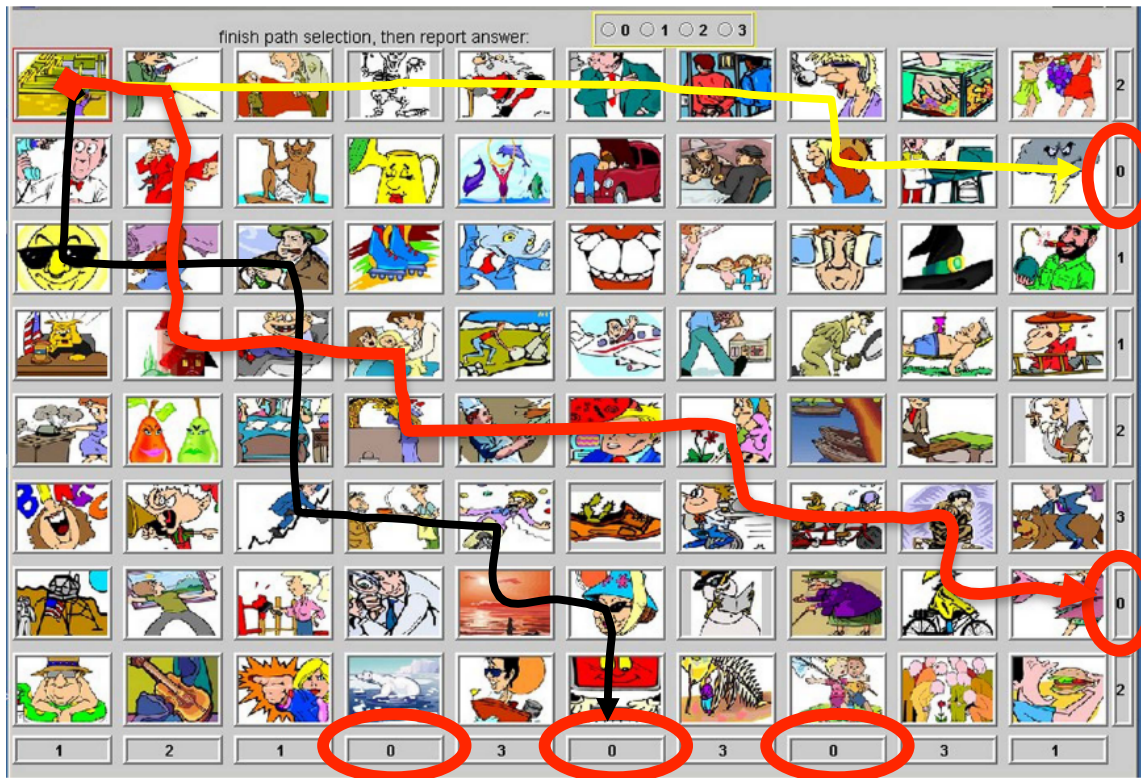


# Score mechanism of probabilistic decision tree

- Rationale:
  - **At least one** of the consistent decision paths is the **correct** path.
  - Other consistent decision paths are “**noises**” whose effects will **cancel out** over multiple rounds.
- Basic probabilities:
  - $P_1$ :  $P(e \in \text{secret}) = k/n$
  - $P_0$ :  $P(e \text{ not } \in \text{secret}) = 1 - P_1$
- Create a *1-element score table*; in each round, compute
  - $P(X) = P(\langle S_1, D_1, D_2, S_2 \rangle) = P_1 * P_0 * P_0 * P_1$
  - $P_C = \text{sum of probabilities of all consistent paths}$
  - $\text{Score}(S_1) += P(X)/P_C$
  - $\text{Score}(D_1) -= P(X)/P_C$

# Statistical bias in decision paths (2/2)

CAS High [S&P06, Weinshall]



43758 possible decision paths in total, with average path length of 14.55.

Secret images score significantly higher than decoy images after a sufficient number of observations.

Recover the exact root secret after observing 65 sessions.

# Usability costs of preventing the two generic attacks

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1. **Large** root secret space

Memory

2. **Large** round secret space

Memory

Computation

3. **Uniformly** distributed challenges

Round  
Number

Window Size

4. **Complex** challenges

Computation

or **counting-based** challenges

Round Number

# Quantitative evidences from psychology

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- Human beings have limitations on **cognitive capability** and **memory**.
  - These limitations **will NOT be** significantly improved even after **repetitive rehearsal**.
- Atomic Cognitive Operations
  - (Single/Parallel) **Recognition**
  - (Free/Cued) **Recall**
  - (Single-target/Multi-target) **Visual Search**
  - Simple **Cognitive Arithmetic**

# High security at cost of heavy cognitive demand

*More secure* ↑

	k	n	Win size	Password space	Reported Time /round(sec)	HP (C) /round (sec)	HP (C) /login (sec)	HP (M)	HP Total =M×C (×10 <sup>2</sup> )
LPN[15]	15	200	200	$1.463 \times 10^{22}$	23.71	33.423	668.45	50.68	338.74
APW[2]	16	200	200	$8.369 \times 10^{24}$	35.50	57.928	347.57	54.05	187.87
CAS Low[31]	60	240	20	$2.433 \times 10^{57}$	5.00	6.073	121.46	70.75	85.94
CAS High[31]	30	80	80	$8.871 \times 10^{21}$	20.00	22.099	220.99	35.38	78.18
SecHCI[20]	14	140	30	$6.510 \times 10^{18}$	9.00	10.638	212.76	16.51	35.13
CHC[32]	5	112	83	$1.341 \times 10^8$	10.97	9.326	93.26	16.89	15.75
PAS[4]	4	N/A	13	$4.225 \times 10^5$	8.37	6.837	68.37	13.51	9.24

↓ *More usable*

The strict tradeoff relation may not hold, but the low bound does.

# Why so hard? – capability asymmetry

## The adversary



### Advantage:

Computation Power  
Storage

### Disadvantage:

Don't know the password

## The user



### Advantage:

Knowledge of the password

### Disadvantage:

Limited cognitive computation

*impossible to do CPA secure encryption*

*$E(\text{secret}, \text{challenge})$*

Limited memory

# Conclusion

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- Our work analyzed the **inherent** limitations of designing Leakage-Resilient Password Systems.
  - Analyze the impact of two **generic attacks** that are usually overlooked.
  - Propose the **design principles** that are necessary to mitigate these generic attacks.
  - Establish the first **quantitative analysis framework** on usability costs of the existing LRPS systems.
- *Our results imply that:*
  - *An LRPS has to incorporate certain **trusted** devices in order to be both **secure** and **usable**.*



Thank You!

**Q & A**



# Brute force for biased challenges

Undercover [CHI08, Sasamoto et al.]: User selects  $k = 5$  pictures out of  $n = 28$ ; # of candidate root secrets is  $C_{28}^5 = 98280$

At most **one** secret image will appear in each challenge.



$P = 1$  (0-indexed)

Brute force recovers the exact root secret after observing 6 sessions.

**Answer** =  $(P + r) \bmod 5$ , where  $r$  is a random integer delivered via a secure channel. Without knowing  $r$ , the answer tells nothing.

# Brute force for round secrets

PAS [ACSAC08, Bai et al.]

(1,1) DFGHKR	(1,2) ABDFGL	(1,3) ABFGJKL	(1,4) DGHLMN	(1,5) CDEFKM
TUVWXYZ	MORSUWY	NSUWXZ	PRUVWXZ	OPSTUXZ
(2,1) DEFHJK	(2,2) CHKLNO	(2,3) AEFHJKQ	(2,4) DEFHJK	(2,5) ABCDEF
OPSTUVW	RSUWXYZ	OPQSTVYZ	GKLMORX	
(3,1) AFGHJK	(3,2) AEFHJKQ	(3,3) BCFHJL	(3,4) AEGHJL	(3,5) DFGHKM
MOQRSTV	RSUWXYZ	OPQWZ	MOQTUVW	NOQTWXY
(4,1) ABFGJK	(4,2) AGHJKM	(4,3) ABCDGH	(4,4) ACEGLM	
NPSTXZ	MQSTUXY	NPQTUWY	LMNOPVX	NPRSTVZ
(5,1) ACEGKM	(5,2) CDEFGH	(5,3) BCHKMN	(5,4) CDEFHJL	(5,5) EFG
NORTWXY	JMOQSTU	RTVWXYZ	MQRSTV	OQRS
(1,1) CEHKLM	(1,2) CEKLNQ	(1,3) ABEGKL	(1,4) ACFLMO	(1,5) ABC
NPQRUVW	PQRSVYZ	OQSTVWY	PQRSUVZ	ORST
(2,1) BCEFMO	(2,2) ACDEFJL	(2,3) ABEJLNQ	(2,4) ACDGHJ	(2,5) ACE
PQSTVWY	NPQTUYZ	KLNQSTX	NQRT	
(3,1) BCDFHJ	(3,2) ADEFGH	(3,3) ABEJLNQ	(3,4) ADEGKM	(3,5) ACD
MNQRSVY	LMPQRUY	RSVWXY	NOQRTU	MOQI
(4,1) BDEKOP	(4,2) ABFGKO	(4,3) ABDEJKL	(4,4) BGE	
QSTUVXZ	NPRSTVW	QSTVWXZ	PSTUVX	OQRSVWX
(5,1) BCDEF LN	(5,2) CDJKNO	(5,3) ABCHKO	(5,4) ACFGJLN	(5,5) ADFHJK
PQRUVX	PQSUXYZ	PRSTVYZ	QRTUVW	NPRVWXZ

Password = ( $\langle 3, 2 \rangle$ , hello),  
 ( $\langle 1, 3 \rangle$ , world)

Challenge index = 2

Predicate = ( $\langle 3, 2 \rangle$ , e), ( $\langle 1, 3 \rangle$ , o)

	2: No No	2: No Yes	2: Yes No	2: Yes Yes
1: No No	WX	RJ	YF	RM
1: No Yes	RJ	RM	WX	YF
1: Yes No	RM	YF	RJ	WX
1: Yes Yes	YF	WX	RM	RJ

Answer = MX

= table[YES, YES][NO, YES]

The SAME index is used for the same authentication session.  
 Brute force recover the round secret after observing 1 session.

Implications: A challenge that can be solved by a small number secret elements is not secure, cognitive workload has to be increased.

# Statistical bias in challenges

Undercover [CHI08, Sasamoto et al.]

At most **one** secret image will appear in each challenge.



$P = 1$  (0-indexed)

Build a 2-element counting table. A secret image will NOT appear together with another secret image. Recover root secret in 20 sessions.

**Answer** =  $(P + r) \bmod 5$ , where  $r$  is a random integer delivered via a secure channel. Without knowing  $r$ , the answer tells nothing.

Implications: A challenge that uniformly draws the candidate elements will be secure, but it will increase the round number or impose a larger window size.

# Statistical bias in responses (1/2)

SecHCI [Cryptography ePrint 05, Li et al.]: Root secret consists of  $k=14$  icons,  $n=140$

Assume the number of your pass-pictures in the following 30 pictures is  $N$ , please tell me  $(N \bmod 4)$  is 0/1 or 2/3.  
Challenge 1:  $N \bmod 4 =$   0/1 or  2/3

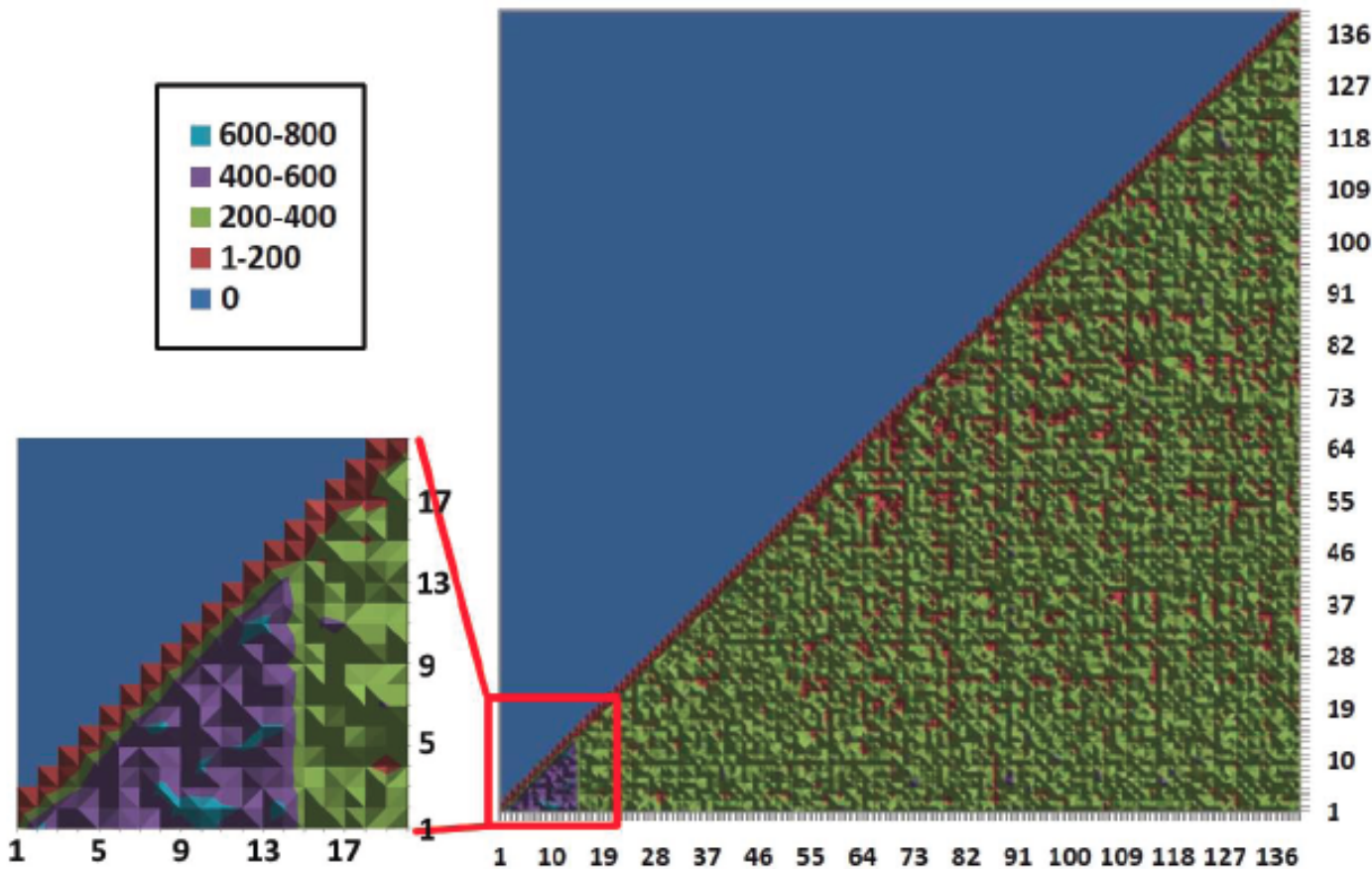


5 secret icons appear (**duplications** allowed).

**Answer** =  $5 \bmod 4 = 1$ , choose '0/1'

Challenges are designed such that 0/1 and 2/3 appear with the same probability, which is different from of the uniform distribution of secret pictures appearing in a challenge.

# Statistical bias in responses (2/2)



Implications: A challenge based on counting problem must use the form  $r = x \bmod 2$ ; otherwise the pair-wise bias appears. This is true for all counting based challenges.

## Usability score in the quantitative analysis framework

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- Cognitive workload:  $HP(C)$ 
  - Measured by sum of the **reaction time** of each atomic operations (e. g., counting, mod, simple arithmetic)
  - How fast can an **average** human solve the challenge?
    - The time limit is implementation-independent
- Memory demand:  $HP(M)$ 
  - Measured by # of elements memorized  $\times$  difficulty factor of the specific memory retrieval operation
    - **Recall** is much more difficult than **recognition**
- $HP = HP(C) \times HP(M)$