

An Interface Specification Language for Automatically Analyzing Cryptographic Protocols

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Overview of Talk

- **Problem and basic solution idea**
 - Cryptographic protocols and network model
 - Protocol failure, TMN example
 - Belief logics
- **Examples of what the Automatic Authentication Protocol Analyzer (AAPA) can do**
 - Interface Specification Language (ISL) TMN specification
 - Terminal output, failed-goals files, other files
 - High points of (unnamed) commercial application
- **Plans for the future**
 - Thorough, but still fast and automatic, analyses
 - Good-guy deductions vrs. Bad-guy searches
- **Lessons for protocol users and designers**



Cryptographic Protocols and Network Model

- **Goal: Secure communication over insecure networks**
 - Networks, principals, and messages
 - No other communications
 - Worst case: enemy can read, or be true source, of everything
 - Confidentiality and authentication
- **Tools for achieving goal**
 - Shared or confirmable secrets
 - Symmetric- and public-key encryption
 - Effectively 1-to-1 hash functions
 - Timestamps, nonces, signatures, etc.
- **Protocols**
 - Distributed algorithms carried out by stages
 - Abort if something not as expected



Protocol Failure

- **Example: TMN (Tatebayeshi-Matsuzaki-Newman) key-distribution protocol**
 - 1. $A \rightarrow S: A, S, B, \{SkA\}_{Rsa(PkS)}$
 - 2. $S \rightarrow B: S, B, A$
 - 3. $B \rightarrow S: B, S, A, \{SkB\}_{Rsa(PkS)}$
 - 4. $S \rightarrow A: S, A, B, \{SkB\}_{Xor(SkA)}$
- **ISL notation, but more-or-less standard**
- **Published (CRYPTO '89), recommended by experts**
- **It's wrong, and has lots of company**

Belief Logics

- **Formalize authentication reasoning that assumes**
 - Good encryption and hash functions
 - Correct distributions of secrets
- **Sample deduction**
 - If P believes only P and Q know K, and P receives an M that K decrypts to something meaningful, then P believes Q sent M -- though not necessarily recently or to P
- **AAPA's BGNY logic**
 - Derived from Gong-Needham-Yahalom (GNY) logic
 - Sending, receiving, belief, freshness, conveyence, shared secrets, possession, recognizability, trustworthiness, not-from-here checks, message extensions, feasibility constraints
 - Many extensions and corrections to GNY (e.g., stages, key-exchange functions, hash codes as keys)

Sample AAPA Analysis: ISL (1)

/* Tatebayashi, Matsuzaki, Newman (TMN) Protocol */

DEFINITIONS:

PRINCIPALS: A,B,S;

PRIVATE KEYS: \wedge PkS;

PUBLIC KEYS: PkS;

SYMMETRIC KEYS: SkA,SkB;

ENCRYPT FUNCTIONS: Xor,Rsa;

Xor WITH ANYKEY HASINVERSE

Xor WITH ANYKEY;

Rsa WITH PkS HASINVERSE Rsa WITH \wedge PkS;

Sample AAPA Analysis: ISL (2)

INITIALCONDITIONS:

A Received Xor,Rsa,A,B,S,PkS,SkA;

**A Believes (Fresh SkA; PublicKey S Rsa PkS;
SharedSecret A S SkA;
Trustworthy B; Trustworthy S);**

B Received Xor,Rsa,A,B,S,PkS,SkB;

**B Believes (Fresh SkB; PublicKey S Rsa PkS;
SharedSecret A B SkB;
Trustworthy A; Trustworthy S);**

S Received Xor,Rsa,^PkS;

**S Believes (PrivateKey S Rsa ^PkS; Trustworthy A;
Trustworthy B);**

Sample AAPA Analysis: ISL (3)

PROTOCOL:

1. **A->S : A,S,B,{SkA}Rsa(PkS);**
2. **S ->B : S,B,A;**
3. **B ->S : B,S,A,{SkB}Rsa(PkS)||(SharedSecret A B SkB);**
4. **S ->A : S,A,B,{SkB}Xor(SkA)||(SharedSecret A B SkB);**

GOALS:

1. **S Possesses SkA;**
3. **S Possesses SkB;**
S Believes SharedSecret A B SkB;
4. **A Possesses SkB;**
A Believes SharedSecret A B SkB;



Sample AAPA Analysis: Terminal Output

Creating theory tmn

Beginning tmn proofs

Initializing globals

Proving default goals, stage 1

Retrying failed default goals, stage 1

Proving user goals, stage 1

Proving default goals, stage 2

Proving user goals, stage 2

Proving default goals, stage 3

Retrying failed default goals, stage 3

Proving user goals, stage 3

User-goal failure, stage: 3!

Goal statement: S Believes (SharedSecret A B SkB);

Sample AAPA Analysis: .fail file (1)

```
/* ##### Failed default goal from stage 3: ##### */  
S Believes  
  (B Conveyed {SkB}Rsa(PkS)||(SharedSecret A B SkB));  
/* ===== Unproved subgoals: ===== */  
S Believes (S Recognizes SkB);  
S Believes (SharedSecret S B SkB);  
S Believes  
  (Fresh {SkB}Rsa(PkS)||(SharedSecret A B SkB));  
/* ===== Proved subgoals: ===== */  
S Received {SkB}Rsa(PkS)||(SharedSecret A B SkB);  
S Possesses Rsa,UNPkS;  
S Believes (PrivateKey S Rsa UNPkS);
```

Sample AAPA Analysis: .fail file (2)

```
/* ##### Failed default goal from stage 1: ##### */  
S Believes (A Conveyed {SkA}Rsa(PkS));  
/* ===== Unproved subgoals: ===== */  
S Believes (S Recognizes SkA);  
S Believes (SharedSecret S A SkA);  
S Believes  
    (Fresh {SkA}Rsa(PkS)||(SharedSecret A S SkA));  
/* ===== Proved subgoals: ===== */  
S Received {SkA}Rsa(PkS);  
S Possesses Rsa,UNPkS;  
S Believes (PrivateKey S Rsa UNPkS);
```

Sample AAPA Analysis: Other Files

- **Have a .thms file giving ISL versions of all theorems**
 - In TMN case, all interesting theorems are proved subgoals
 - In other cases, useful for figuring out what happened
- **Have option of producing**
 - Higher Order Logic (HOL) theory of protocol
 - HOL translation of ISL input
- **Optional outputs mainly used for debugging AAPA**

AAPA Analysis of Commercial Protocols

- **Customer requested confidentiality**
- **Protocols moderately complicated, and huge**
 - Roughly 100 items or subitems in some messages
 - Most of detail irrelevant to AAPA analysis -- but which?
 - Biggest formally analyzed examples known to author
- **Results of analyses of two protocols**
 - Did not find failures in protocols
 - Found omissions, errors and inconsistencies in documentation
 - Produced basis for much better documentation
- **AAPA additions necessary for effort**
 - ISL abbreviation capacity ($X = Y$ as in C)
 - New diagnostics for feasibility failures

Plans for the Future

- **Belief logics vs. attack construction**
 - **Simplicity and speed vs. thoroughness and rigor**
 - **How to gain one without losing the other: Replace searches with using theorems about searches**
- **Research program**
 - **Find failed protocols in literature**
 - **Analyze them with AAPA**
 - **For failures AAPA misses, ask where belief logic allowed false beliefs during attacks**
 - **Adjust logic and repeat process; time always polynomial**
- **User interfaces**
 - **Use CAPSL (Common Authentication Protocol Specification Language) by Millen and protocol-analysis community**
 - **Make sure ISL virtues survive in CAPSL**

Lessons for Protocol Users and Designers

- **Protocol failure is a little-known, but very real problem**
 - “Easy” design problem is actually a weak link
 - About half of published protocols fail -- estimate based on Lowe’s and my own experiences
- **An AAPA analysis is worth performing**
 - Finds common failures
 - Gives overview, corrects documentation, identifies information flows, and identifies trust assumptions
 - It’s fast and cheap
- **A near-future AAPA could make protocol failure, for practical purposes, into a solved problem**