Pitfalls in Designing Zero-Effort Deauthentication: Opportunistic Human Observation Attacks

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The deauthentication problem

- Threat:
 - unauthorized access to a terminal
 - after legitimate user has walked away
- What we actually want is **zero-effort** deauthentication
- Both innocent and malicious adversaries

Zero-effort deauthentication systems

- Already in use!
 - BlueProximity



- Keyless Entry in high end cars
- Based on short-range wireless channels: RSS from user devices





ZEBRA: a recent proposal for deauthentication

Targeted for hospital wards, factory floors, ... User may step away from Terminal but lingers nearby



• No user profiling!

[1] Mare, et al., "ZEBRA: Zero-effort bilateral recurring authentication." IEEE Symposium on Security and Privacy (SP) 2014

http://dx.doi.org/10.1109/SP.2014.51 Mika Juuti: Pitfalls in Designing Zero-effort Deauthentication 4

ZEBRA works by averaging out misclassifications [1]



Window size 10, Threshold 70% 8/10 matches ≥ 70% User remains logged in

Bracelet data \rightarrow classes:

- 1. (any) typing
- 2. (any) scrolling
- 3. mouse \leftrightarrow keyboard movements (MKKM)

Only interactions seen at Terminal considered [1]



– Why? User privacy [1], accuracy of classifier?

ZEBRA vs malicious attackers [1]



- Attacker required to mimic all of victim's interactions
- 20 participants as attackers; researchers as victims
 - Victims verbally announce their interactions

Does ZEBRA resist malicious attackers? [1]



g = deauthentication
at # failed windows

Average window length = 6s

Fraction of adversaries remaining logged in (window size = 21, threshold=60%)

Is this a reasonable adversary model?

More realistic adversary models

- 1. Naïve all-activity
 - As in Mare et al [1]: mimics all
- 2. Opportunistic keyboard-only
 - Mimics selected typing
- 3. Opportunistic all-activity

Mimics selected activities

4. Audio-only opportunistic KB-only – Mimics selected typing, but no line of sight



Our implementation of ZEBRA

• Implemented end-to-end ZEBRA from scratch

- Using off-the-shelf Android Wear smartwatch
 Wider applicability: existing affordable models
- Re-use ZEBRA parameters/methodology wherever possible

Parameter comparison

Parameter name	Original implementation	Our implementation
Minimum duration	25 ms	25 ms
Maximum duration	1 s	1 s
Idle threshold	1 s	1 s
Window size	21	20
Match threshold	60%	60%
Overlap fraction	Not reported	0
Grace period	1, 2	1, 2
Classifier	Random forest	Random forest
Classifier training data	Form filling	Form filling
Validation methodology	Not reported	Leave-one-user-out

• Bracelet hardware, datasets used...

Our implementation Architecture



Android Wear application for smartwatch Matlab Random Forest classifier for interaction classification Java application for Terminal Mika Juuti: Pitfalls in Designin

Our implementation of ZEBRA (2)



Zebra/java\$ find -name *.java -print | xargs grep -v "\\\\" |
grep -v "1\$" | grep -v "*" | wc -l
Zebra/java\$ 7706
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Naïve malicious attackers: comparison

at # failed windows Fraction of adversaries with access in window w 1.0 q=10.9 $\alpha = 2$ g=20.8 0.8 Fraction of logged in adversaries 0.7 0.6 0.6 0.5 0.4 0.4 0.3 0.2 0.2 0.1 0 5 10 15 20 25 30 12 0 2 6 8 10 14 Δ Windows (w) Window (w) Original malicious attacker (naïve) [1] Our naïve all-activity attacker

- 20 participants as victims; researchers as attackers

All attackers are deauthenticated

g = deauthentication

ZEBRA does not resist opportunistic malicious attackers

g = deauthentication at # failed windows



- 20 participants as victims; researchers as attackers
- Attackers do not eventually get logged out

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Can still protect against innocent "attackers"

- mismatched traces model innocent attackers
- All users eventually deauthenticated
- Avg. window length = 14s



g = deauthentication

Mismatched user traces

What went wrong? [1]

- 1. Inadequate **adversary modeling** in [1]!
- 2. Fundamental design flaw in ZEBRA:

"Authentication based on input source controlled by adversary"

- Attacker controls Terminal:
 - Can choose type/timing of interactions
- A case of tainted input:



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Strengthening ZEBRA [1]

- Recognizing more terminal interactions
- Recognizing off-terminal interactions!
- Black/whitelisting, sanitizing input
- Augmenting with trusted input: RSS

Take-home message

- 1. **Zero-effort security** is appealing
 - Balance between usability and security
 - Care in defining adversary model
- 2. ZEBRA susceptible to **opportunistic attackers**, still effective for preventing **accidental misuse**

Ask me for a demo!